



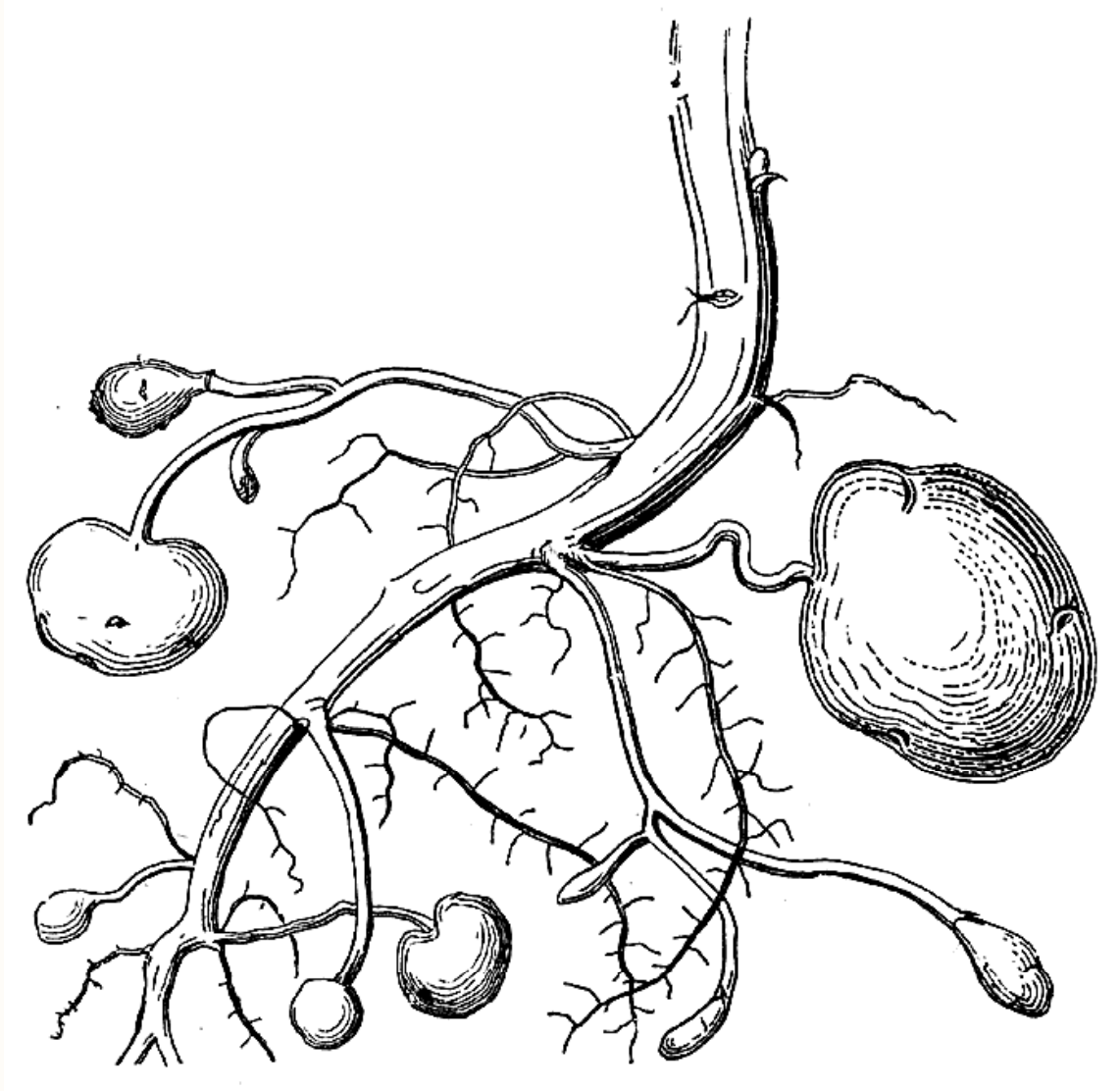
# ROOT

Bertrand Bellenot, Axel Naumann  
CERN

# What's ROOT?



# What's ROOT?



# What's ROOT?



# What's ROOT?



The screenshot displays the ROOT software interface with several windows and plots:

- Top Window:** A browser window showing the ROOT website URL: `http://root.cern.ch/root/html/TParallelCoord.html`.
- Left Panel:** Three histograms showing distributions of `Mbb in (z-dB<0.1) vs Mbb`, `Mbb in/Out`, and `Mbb in - Out`.
- Center Panel:** A candle chart titled `TParallelCoord` showing data points for variables: `Age`, `Grade`, `Step`, `Cost`, `Division`, and `Nation`. The chart shows a dense network of lines connecting data points across these variables.
- Right Panel:** Three histograms with fit statistics:
  - Top Histogram:** `Mbb` vs `Mbb`. Shows a distribution of `Mbb` values with a fit. Statistics: Entries: 67790, Mean: 74.23, RMS: 30.41, Underflow: 0, Overflow: 210.6.
  - Middle Histogram:** `Ratio of Mbb_in/Mbb_out_corr` vs `Mbb`. Shows a ratio distribution with a fit. Statistics: Entries: 44, Mean: 224.3, RMS: 90.93, Underflow: 0, Overflow: 0,  $\chi^2 / \text{ndf}$ : 39.63 / 27, Prob: 0.0005533, p1: 1.055 ± 0.016, p2: 0.0002436 ± 0.0002326.
  - Bottom Histogram:** `Difference of Mbb_in-Mbb_out_corr` vs `Mbb`. Shows a difference distribution with a fit. Statistics: Entries: 3646, Mean: 73.99, RMS: 30.48, Underflow: 0, Overflow: 210.6,  $\chi^2 / \text{ndf}$ : 9.82 / 8, Prob: 0.2779, Constant: 2.835e+004 ± 8251, Mean: 86.98 ± 5.37, Sigma: 9.522 ± 4.888.

# ROOT: An Open Source Project



- Started in 1995
- 11 full time developers at CERN, plus Fermilab, Agilent Tech, Japan, MIT (one each)
- Large number of part-time developers: let users participate
- Available (incl. source) under GNU LGPL



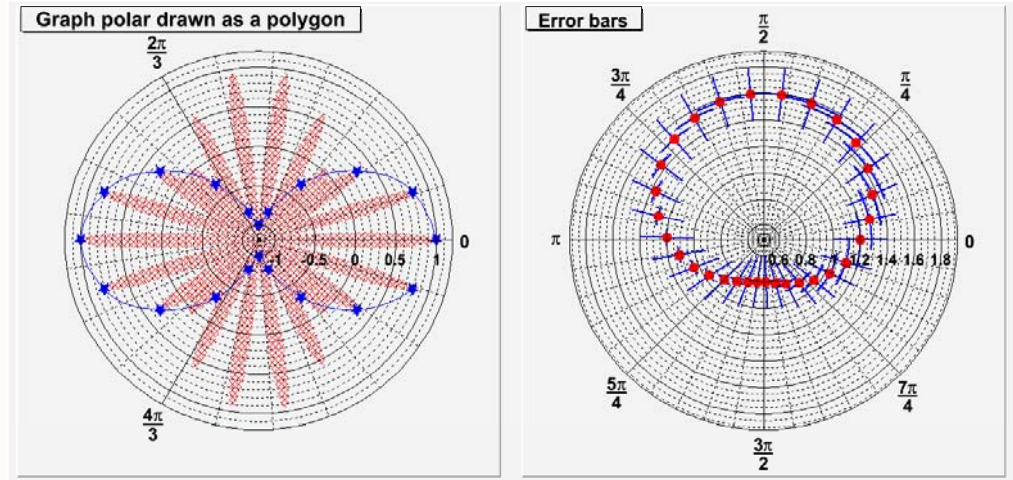
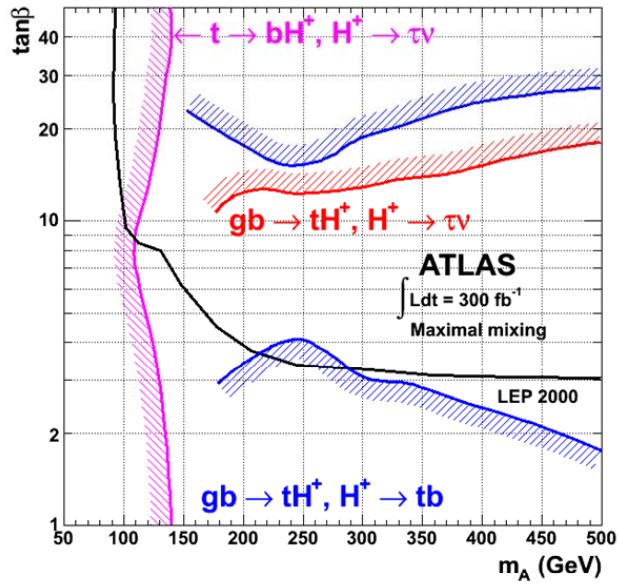
# ROOT in a Nutshell

Framework for large scale data handling

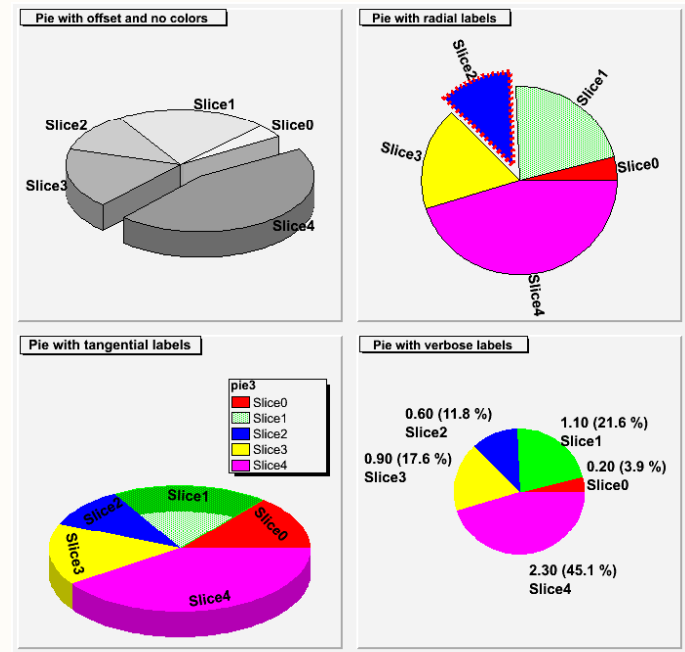
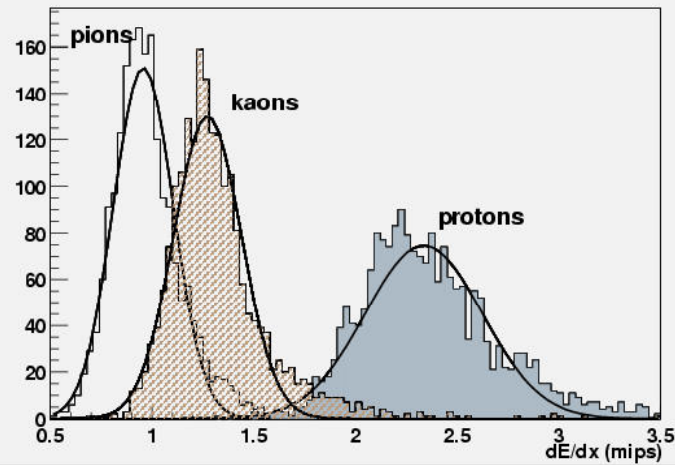
Provides, among others,

- an efficient data storage, access and query system (PetaBytes)
- advanced statistical analysis algorithms (multi dimensional histogramming, fitting, minimization and cluster finding)
- scientific visualization: 2D and 3D graphics, Postscript, PDF, LateX
- geometrical modeller
- PROOF parallel query engine

# Graphics



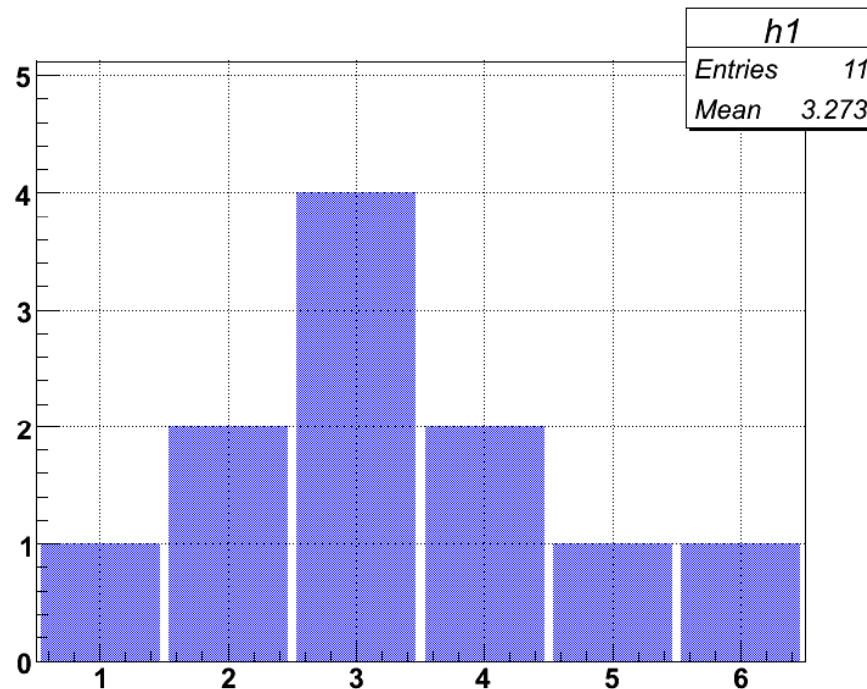
Momentum 730-830 MeV/c





# Histogramming

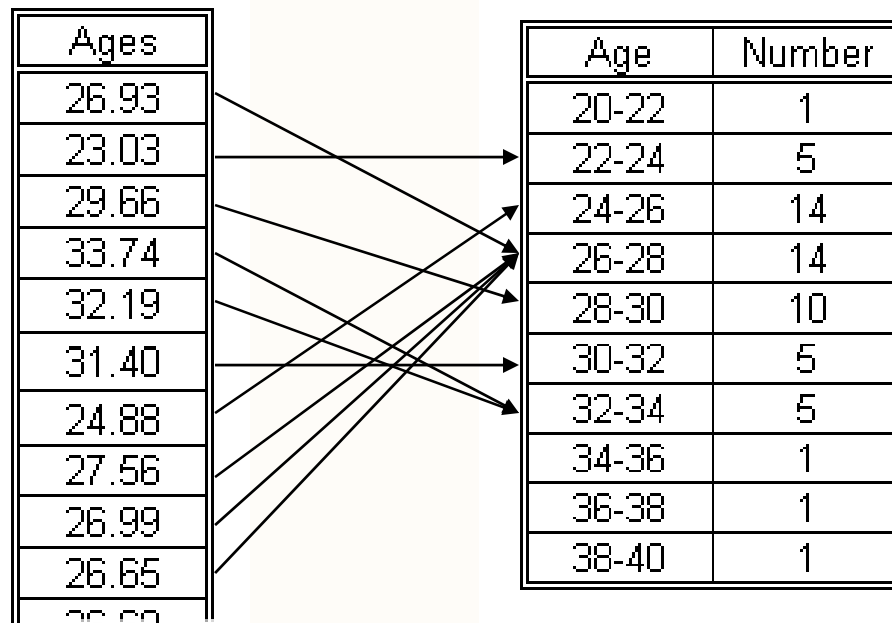
- Histogram is just occurrence counting, i.e. how often they appear
- Example: {1,3,2,6,2,3,4,3,4,3,5}



# Histogramming

- **How is a Real Histogram Made?**

Lets consider the age distribution of the participants:



**Binning:**

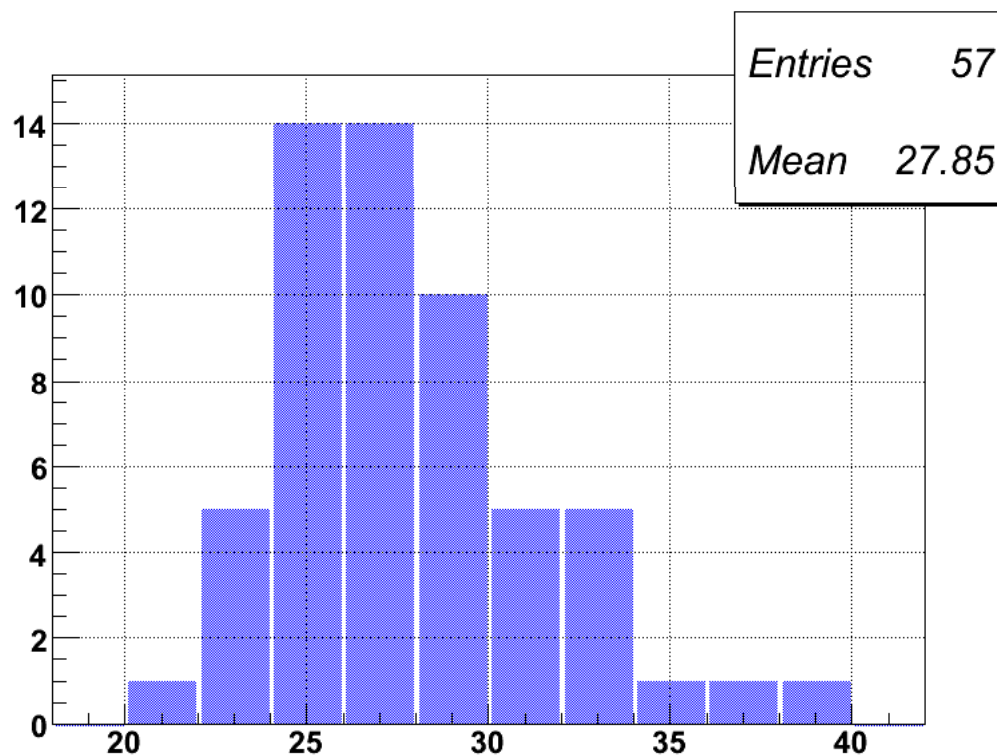
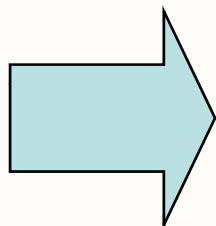
**Grouping ages of participants in several categories (bins)**

# Histogramming



Table of Ages  
(binned)

Age	Number
20-22	1
22-24	5
24-26	14
26-28	14
28-30	10
30-32	5
32-34	5
34-36	1
36-38	1
38-40	1



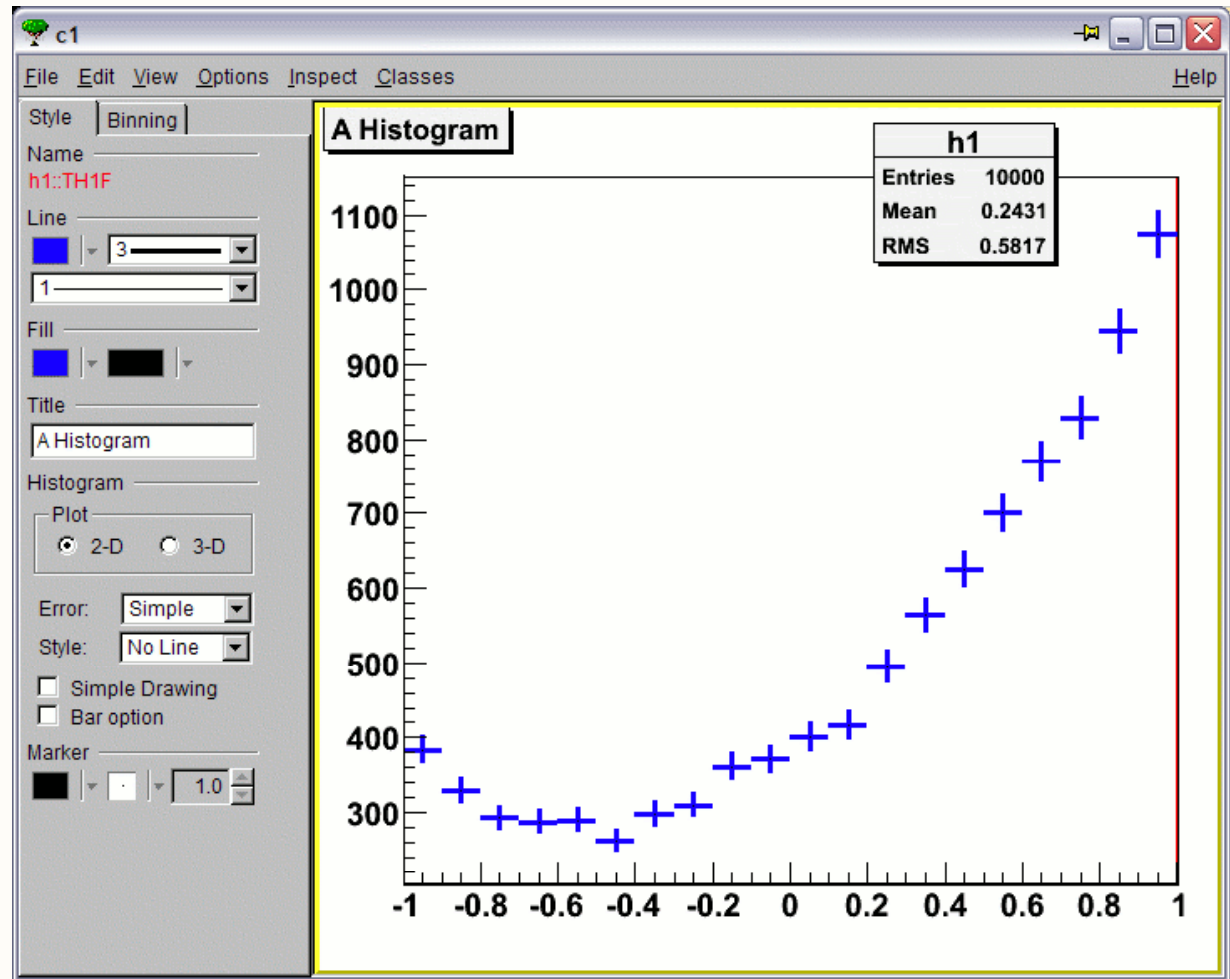
Shows distribution of ages, total number of entries (57 participants) and average: 27 years 10 months 6 days...

# Histograms



Analysis result: often a histogram

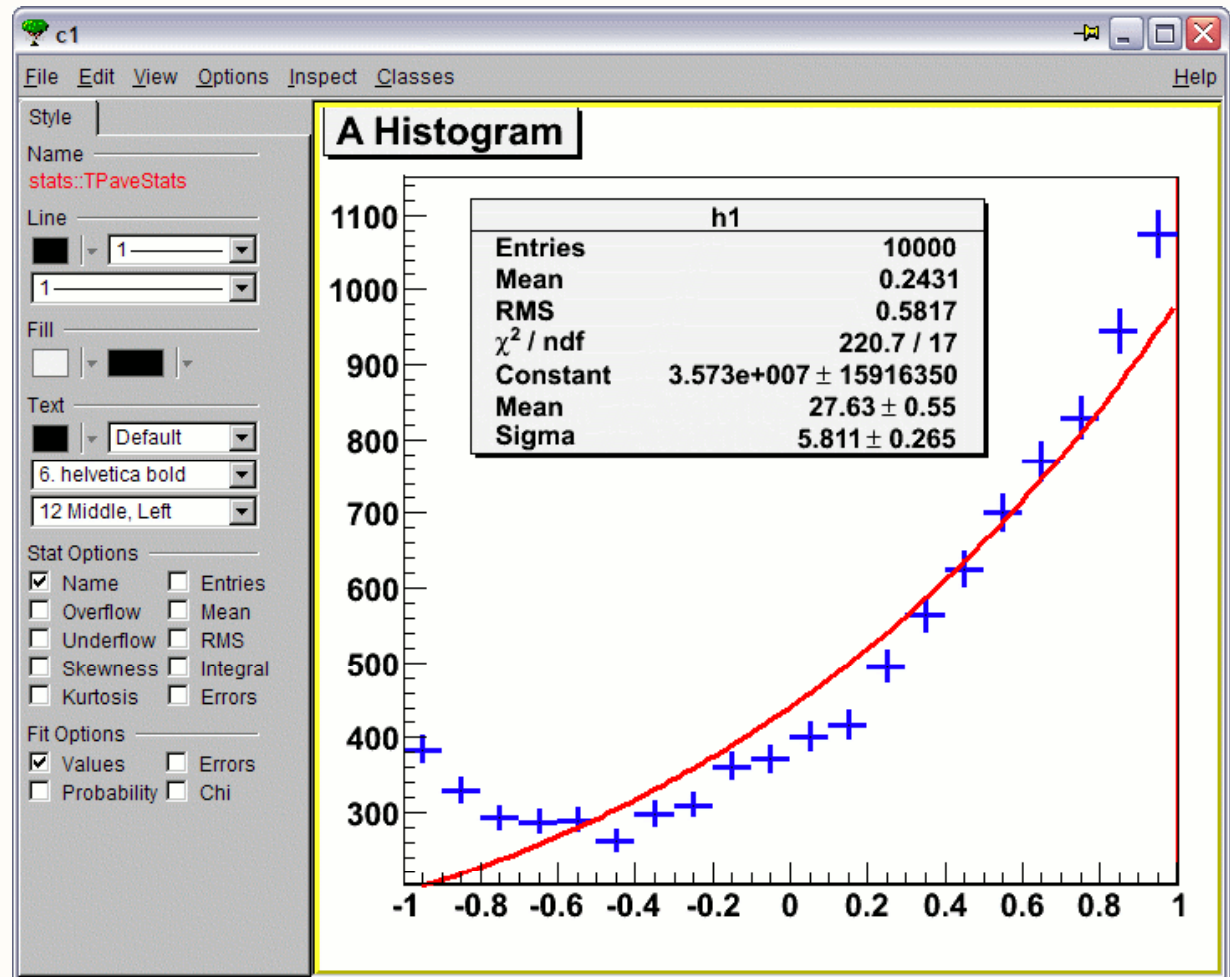
Menu:  
View / Editor



# Fitting



Analysis result:  
often a *fit*  
based on a  
histogram



# Fit



Function describing the distribution of data

Fit = optimization in parameters,

e.g. Gaussian  $f(x) = [0] \cdot e^{-\frac{(x-[1])^2}{2 \cdot [2]^2}}$

For Gaussian: [0] = "Constant"

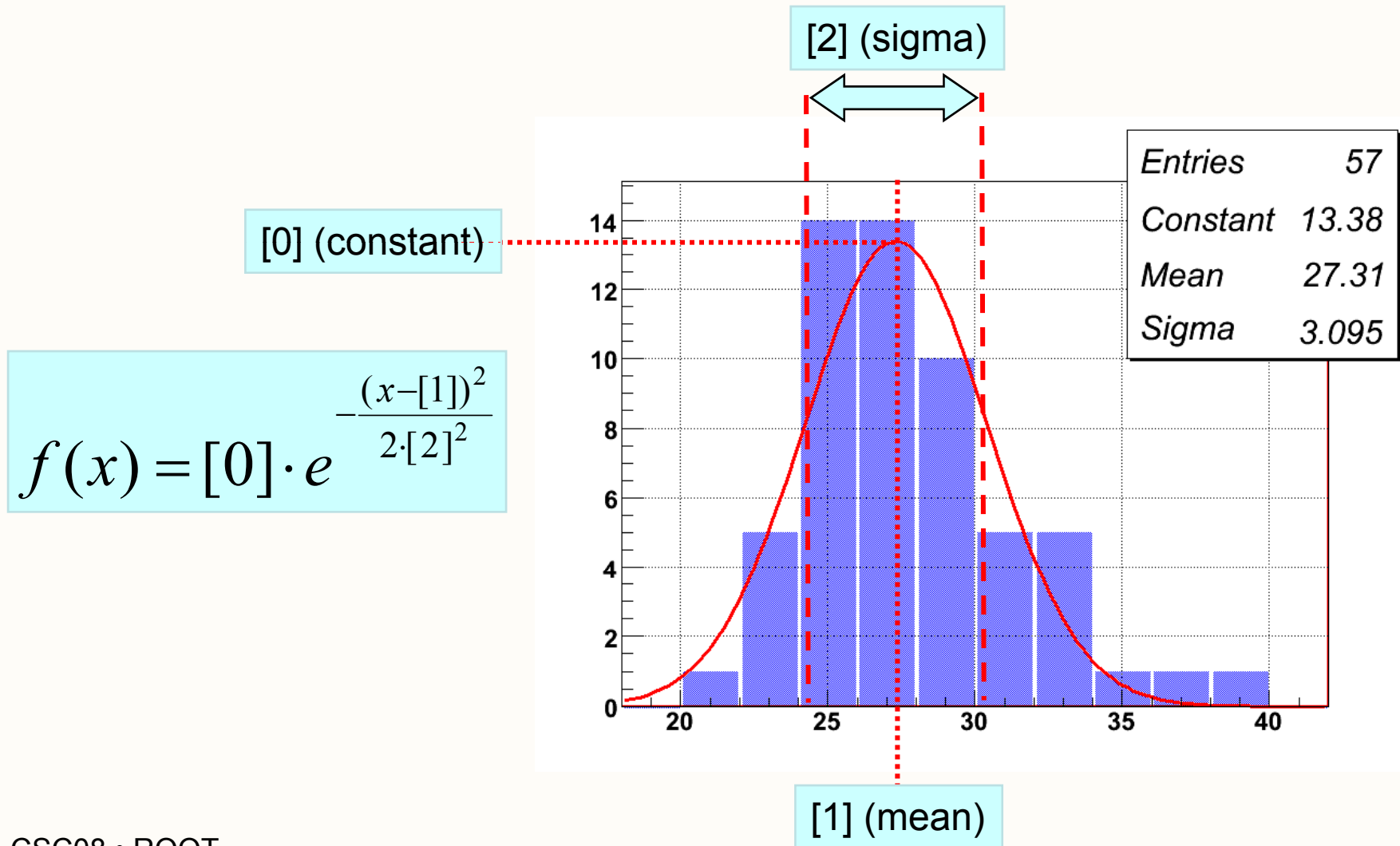
[1] = "Mean"

[2] = "Sigma" / "Width"

Objective: choose parameters [0], [1], [2] to get function as close as possible to histogram

# Fit

Gaussian fit over our histogram:

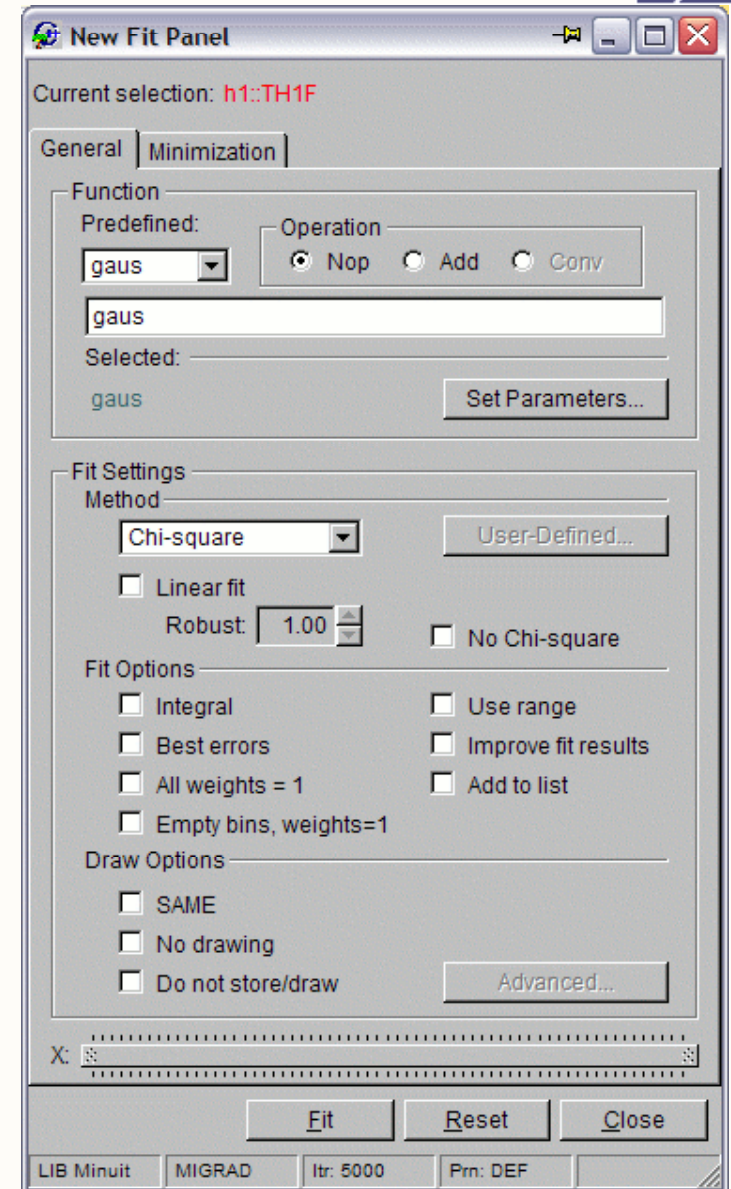


# Fit Panel



To fit a histogram:  
right click histogram,  
"Fit Panel"

Straightforward interface  
for fitting!





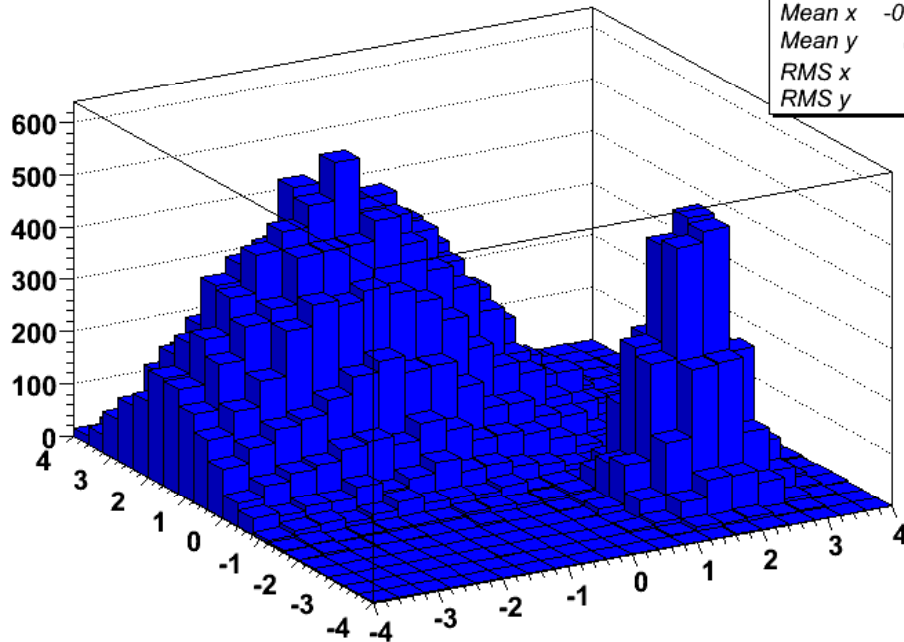
# 2D/3D



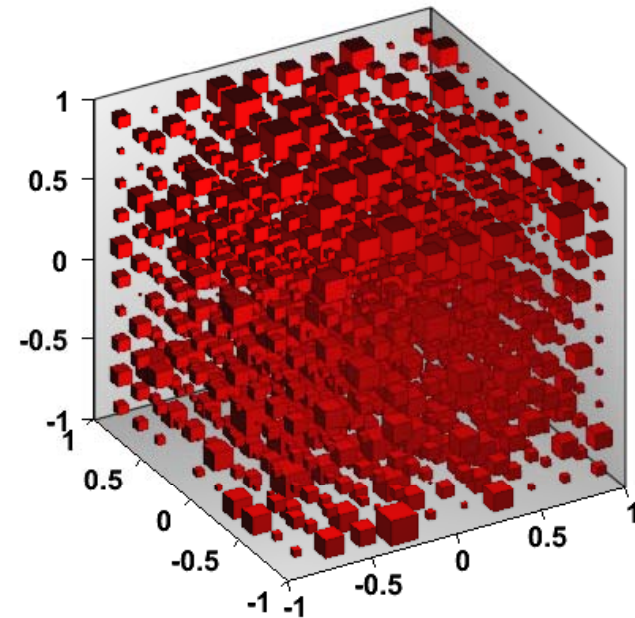
We have seen 1D histograms, but there are also histograms in more dimensions.

```
xygaus + xygaus(5) + xylandau(10)
```

h2	
Entries	40000
Mean x	-0.6685
Mean y	0.967
RMS x	1.826
RMS y	1.541



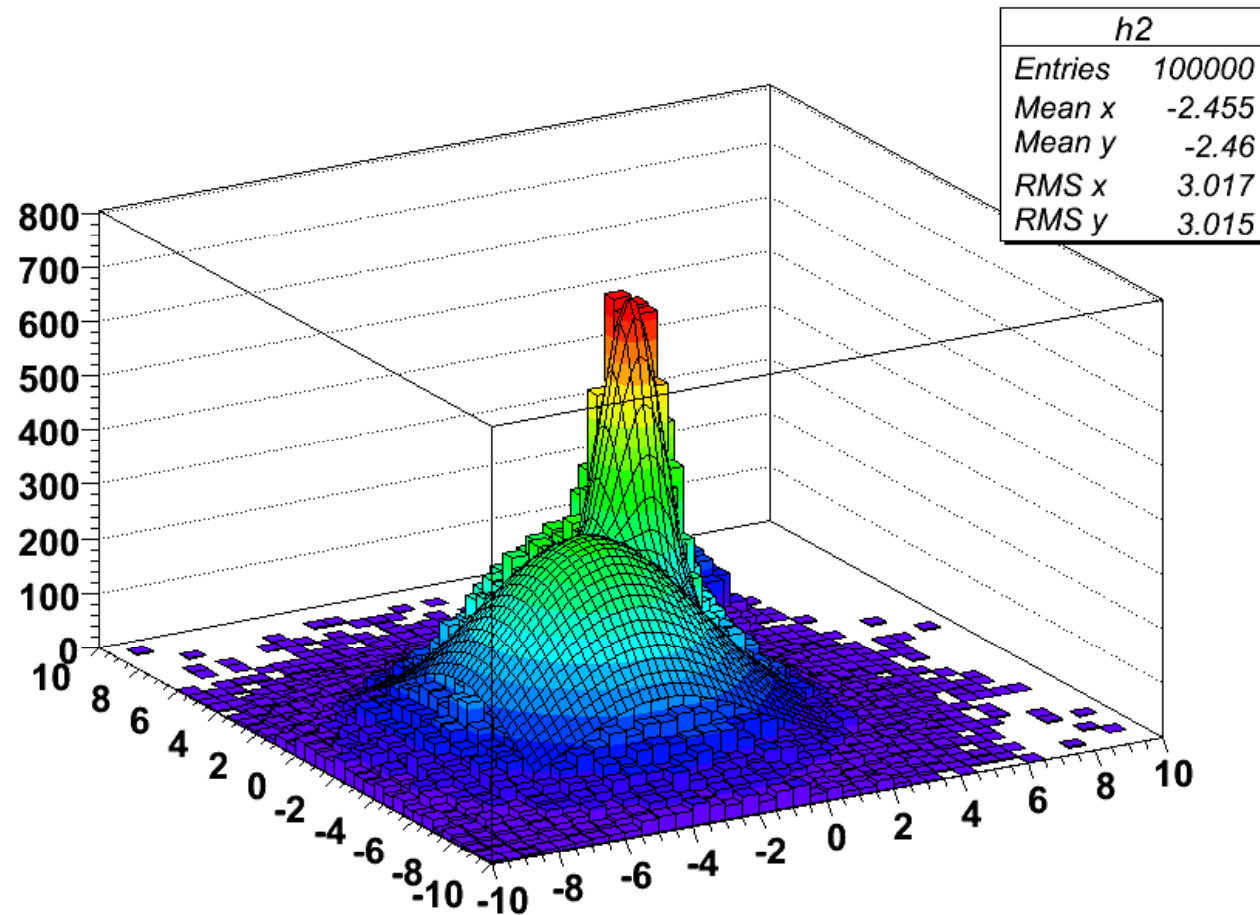
2D Histogram



3D Histogram

# 2D Fitting

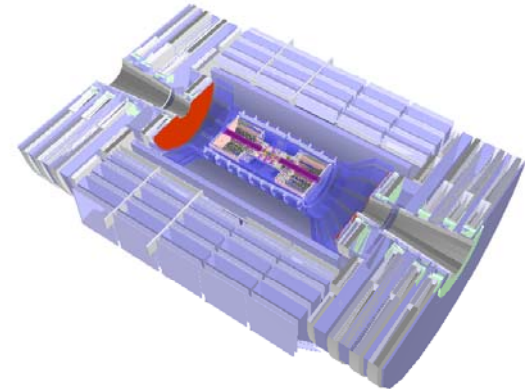
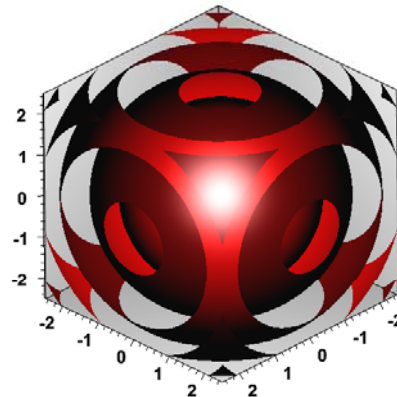
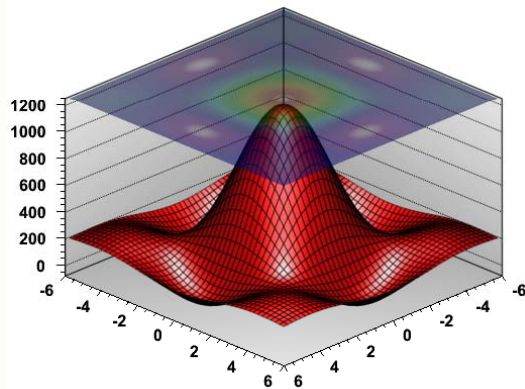
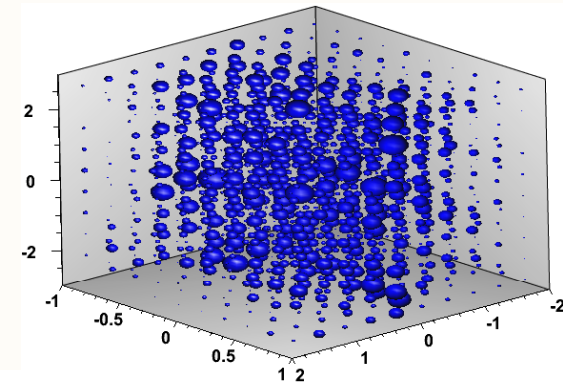
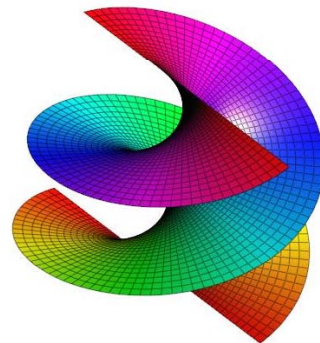
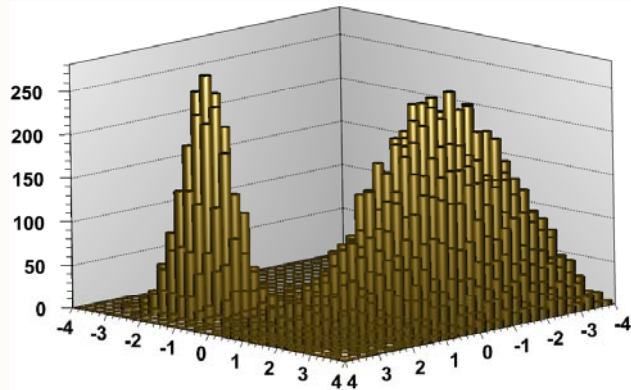
Example of a fit over a 2D histogram



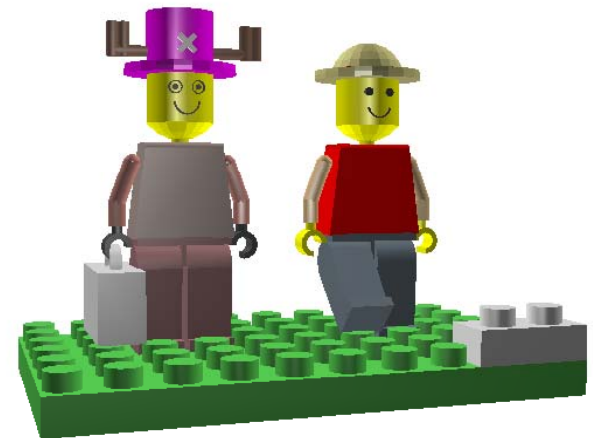
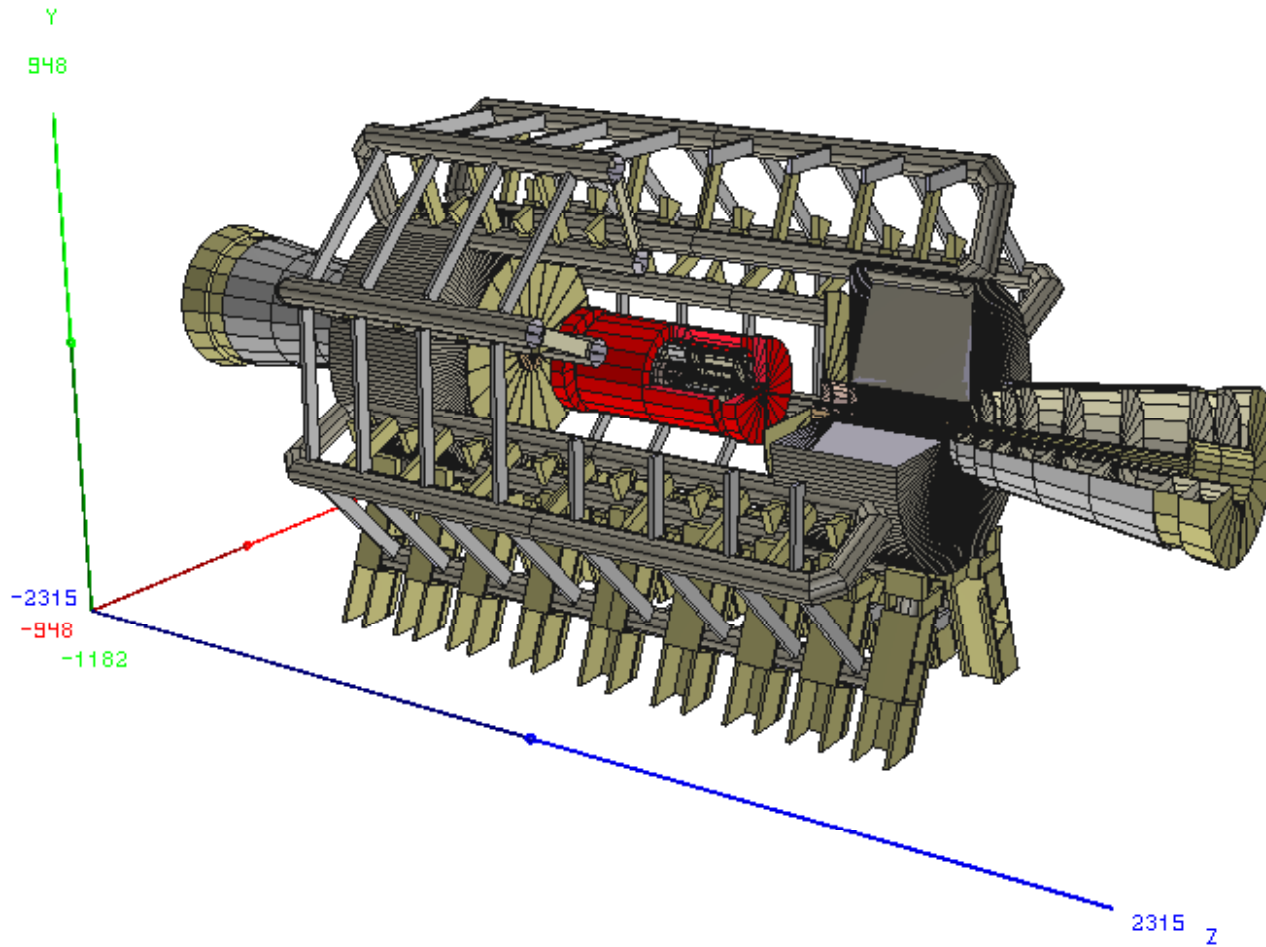
# OpenGL



- OpenGL can be used to render 2D & 3D histograms, functions, parametric equations, and to visualize 3D objects (geometry)



# OpenGL

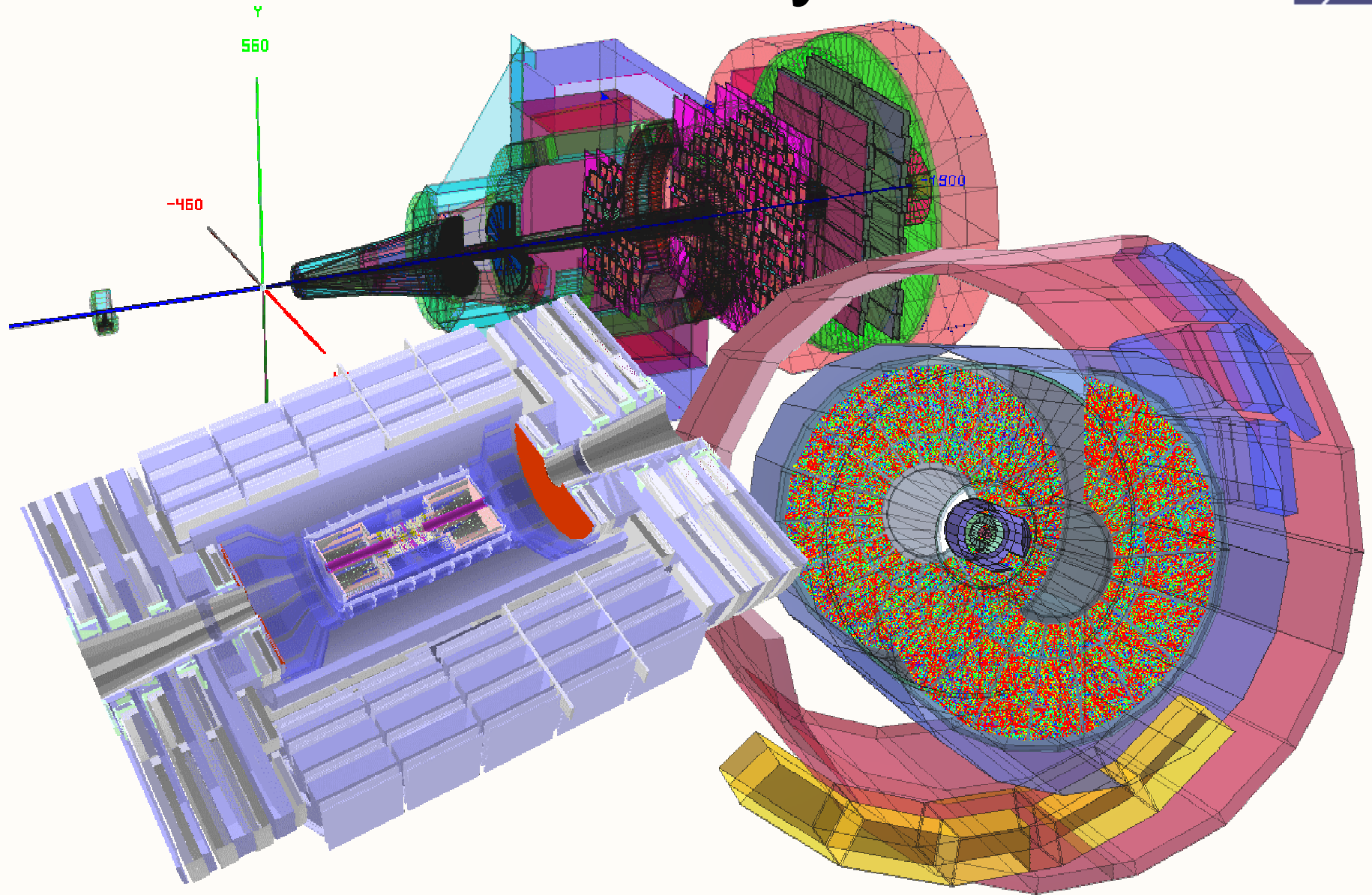


# Geometry



- Describes complex detector geometries
- Allows visualization of these detector geometries with e.g. OpenGL
- Optimized particle transport in complex geometries
- Working in correlation with simulation packages such as GEANT3, GEANT4 and FLUKA

# Geometry



# EVE (Event Visualization Environment)



- Event: Collection of data from a detector (hits, tracks, ...)

Use EVE to:

- Visualize these physics objects together with detector geometry (OpenGL)
- Visually interact with the data, e.g. select a particular track and retrieve its physical properties

# EVE



The screenshot displays the EVE (Event Viewer) software interface. The main window, titled "Eve Main Window", shows a 3D visualization of the ALICE detector components and particle tracks. The left sidebar contains a tree view of the event structure, including EMCAL Hits, ITS Hits, FMD Hits, PMD Hits, TO Hits, TOF Hits, TPC Hits, TRD Hits, VZERO Hits, and ESD Tracks. Below the tree is a style configuration panel for the selected "ESD Tracks" object, showing options for name, TEveElement, show settings, marker, line, and render style.

The central 3D view shows a top-down view of the detector with various colored components and particle tracks. Two smaller 2D projection views are visible at the bottom left: "Rho-Z" and "R-Phi".

Two summary tables are shown, both titled "Alice Event Display Summary Table". The top table lists 10 tracks, and the bottom table lists 6 tracks. Each table has columns for Momentum, P<sub>t</sub>, Phi, Theta, and Eta.

Alice Event Display Summary Table				
ESD Tracks				
Momentum	P <sub>t</sub>	Phi	Theta	Eta
<input type="checkbox"/>	0.5796	0.4517	-2.5215	0.8934
<input type="checkbox"/>	0.4516	0.3285	-2.3953	0.8146
<input type="checkbox"/>	0.5304	0.4083	0.0665	0.8787
<input type="checkbox"/>	0.6510	0.6292	0.0765	1.8305

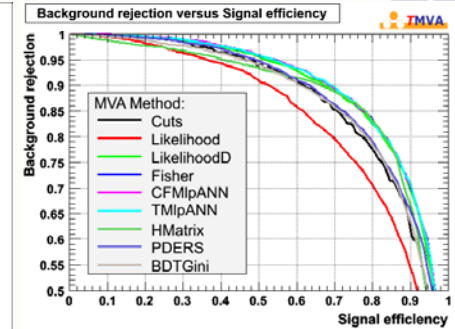
Alice Event Display Summary Table				
ESD Tracks				
Momentum	P <sub>t</sub>	Phi	Theta	Eta
<input type="checkbox"/>	0.8899	0.6894	-1.3460	2.2553
<input type="checkbox"/>	0.3562	0.3514	-0.5208	1.7342
<input type="checkbox"/>	0.2646	0.2084	-1.5928	2.2352
<input type="checkbox"/>	0.2207	0.1725	-1.7871	0.8969



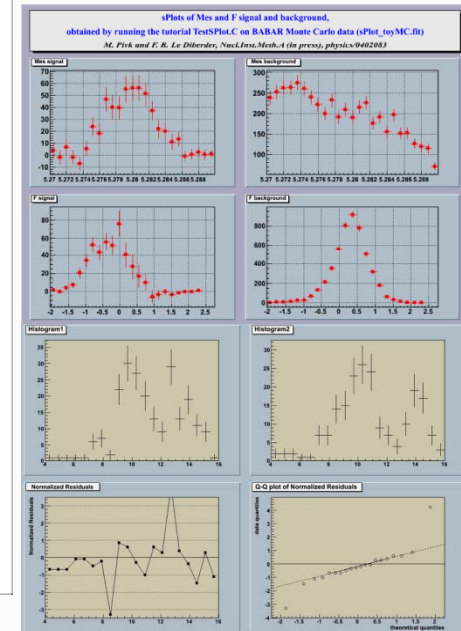
# Math



## TMVA



## SPlot



Histogram library

TH1 TF1

MathMore

Random Numbers

Extra algorithms

Extra Math functions

GSL and more

MathCore

Function interfaces

Physics Vectors

Basic algorithms

Basic Math functions

Statistical Libraries

Statistical Utilities

TMVA MLP

Fitting and Minimization

New Fitter

RooFit

TMinuit

TFumili

Linear Fitter

Minuit2 (new C++ Minuit)

Linear Algebra

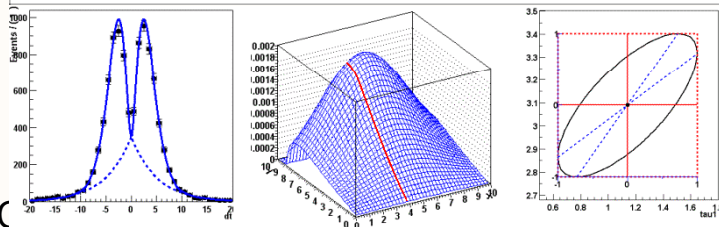
TMatrix

SMatrix

libCore

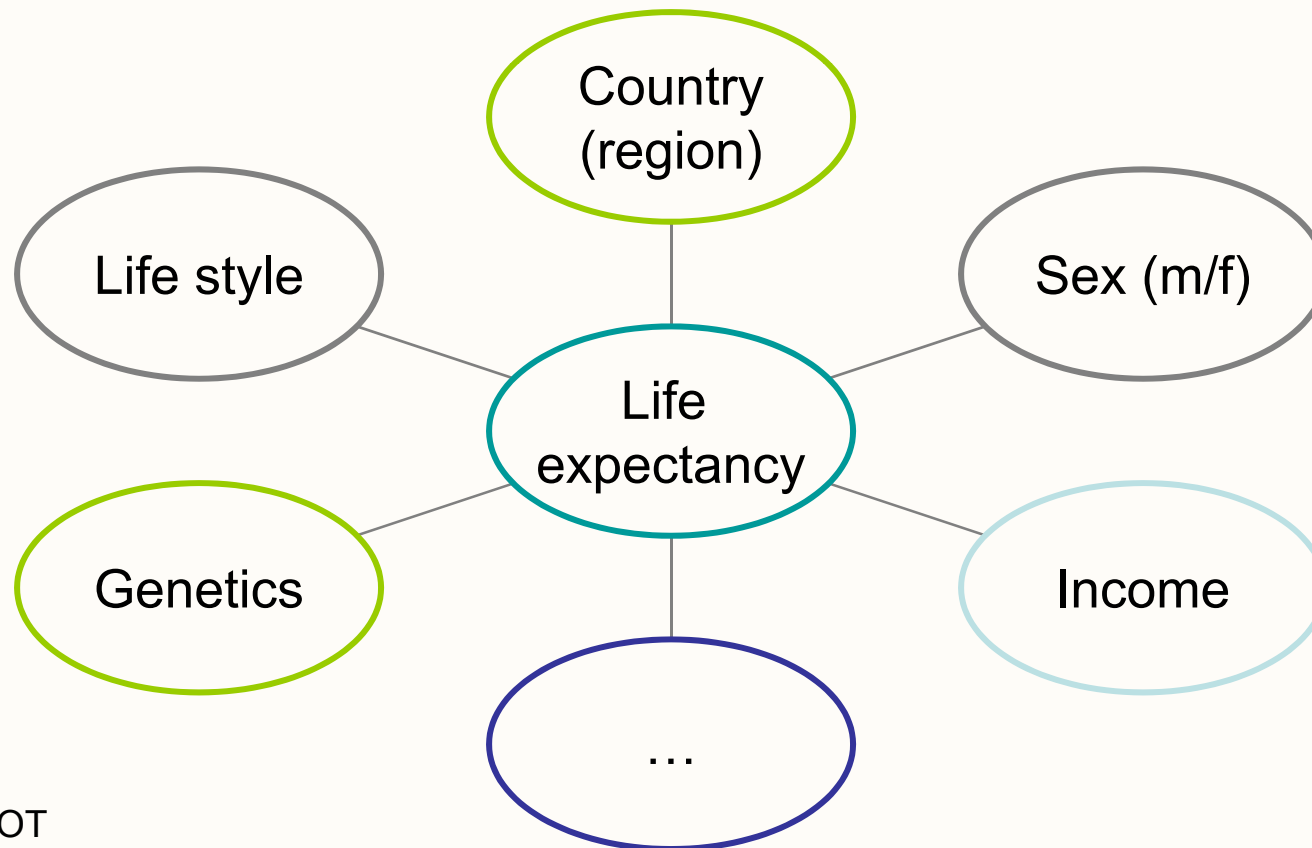
TMath

TRandom



# Multivariate Analysis

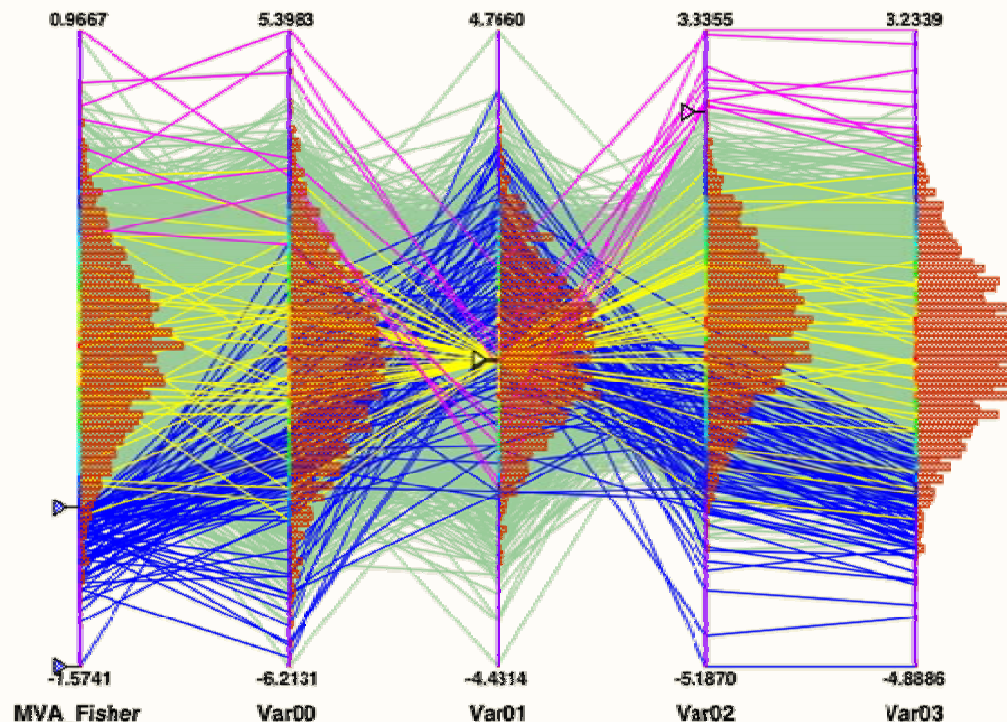
- Consider this simple question: How to estimate someone's life expectancy?
- This depends on many variables:



# Multivariate Analysis



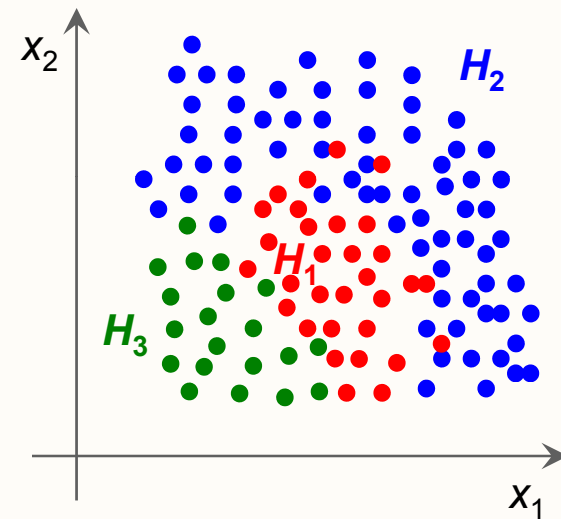
- Many variables? *Parallel Coordinates*



- This will not help to solve the problem, it only allows to visualize multiple variables

# Multivariate Analysis

- Sample described by  $k$  variables (that are found to be discriminating)
- Samples can be classified into  $n$  categories:  $H_1 \dots H_n$
- E.g.
  - $H_1$  : life exp.  $< 40$
  - $H_2$  : life exp.  $40..60$
  - $H_3$  : life exp.  $> 60$
- Example:  $k=2$  variables  $x_1, x_2$   
 $n=3$  categories  $H_1, H_2, H_3$

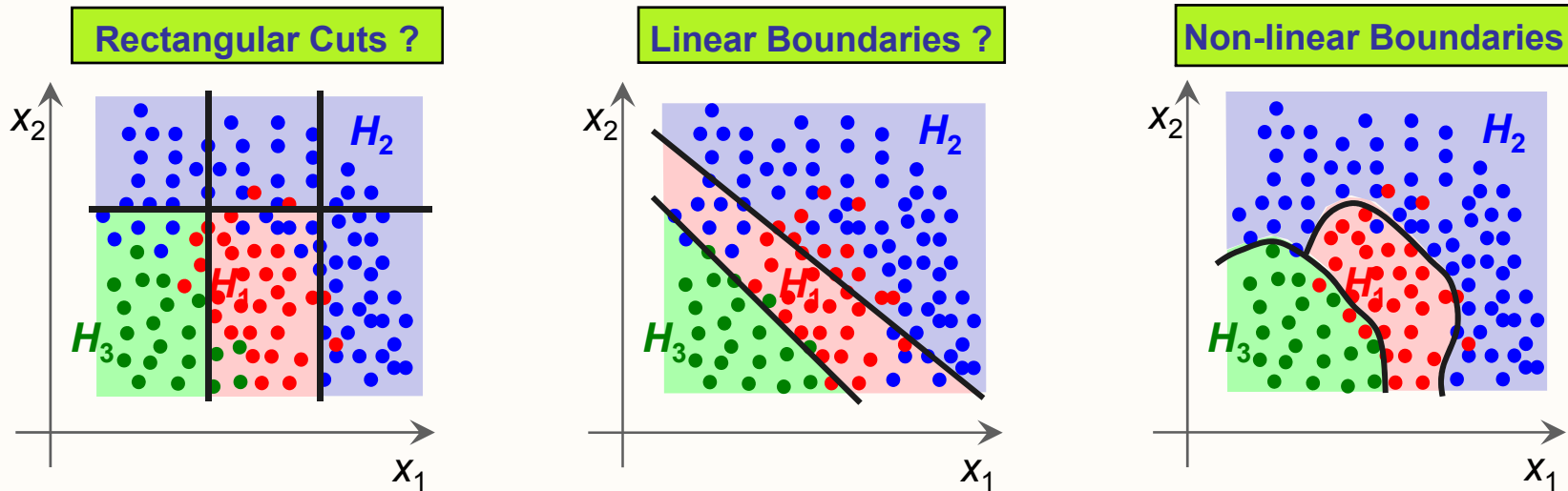


Example:  $k=2, n=3$

# Multivariate Analysis



Problem: Find boundaries between  $H_1$ ,  $H_2$ , and  $H_3$  such that  $f(\mathbf{x})$  returns the category of  $\mathbf{x}$  with maximum correctness



Simple example  $\rightarrow$  I can do it by hand.

Large input variable space, complex correlations: manual optimization very difficult

# Multivariate Analysis



Generic problem: find category for set of values

To make such an estimation, we need two phases:

- (Supervised) learning / training phase:
  - Take samples for which categories are known
  - Machine adapts to give the smallest classification error on training sample
- Processing phase:
  - The trained system can now analyze and produce output for any new sample

# TMVA



- Framework offering a collection of data mining tools, e.g. NN (Neural Network), GA (Genetic Algorithm), ...
- In HEP mostly two class problems – signal (S) and background (B)
  - Physics processes
  - Finding physics objects
  - Detector readout
  - ...

# Interlude: HELP!



ROOT is a framework – only as good as its documentation.

User's Guide (it has your answers!)

<http://root.cern.ch>

“What is TNamed? What functions does it have?”

Reference Guide

<http://root.cern.ch/root/html>



# Let's fire up ROOT!





# Setting Up ROOT

Before starting ROOT:

setup environment variables \$PATH,  
\$LD\_LIBRARY\_PATH

Go to where ROOT is:

```
$ cd /path-to/root
```

```
(ba)sh: $ . bin/thisroot.sh
```

```
(t)csh: $ source bin/thisroot.csh
```

# Starting Up ROOT



ROOT is prompt-based

```
$ root  
root [0] _
```

Prompt speaks C++

```
root [0] gROOT->GetVersion();  
(const char* 0x5ef7e8) "5.16/00"
```

# ROOT As Pocket Calculator



Calculations:

```
root [0] sqrt(42)  
(const double)6.48074069840786038e+00  
root [1] double val = 0.17;  
root [2] sin(val)  
(const double)1.69182349066996029e-01
```

Uses C++ Interpreter CINT



# Running Code

To run function mycode() in file mycode.C:

```
root [0] .x mycode.C
```

Equivalent: load file and run function:

```
root [0] .L mycode.C  
root [1] mycode()
```

Quit:





```
root [0] .q
```

All of CINT's commands (help):

```
root [0] .h
```

# ROOT Prompt



-  Why C++ and not a scripting language?!
-  You'll write your code in C++, too. Support for python, ruby,... exists.
-  Why a prompt instead of a GUI?
-  ROOT is a programming framework, not an office suite. Use GUIs where needed.

# Running Code



Macro: file that is interpreted by CINT (**.x**)

```
int mymacro(int value)
{
    int ret = 42;
    ret += value;
    return ret;
}
```

Execute with **.x mymacro.C(42)**

# Compiling Code: ACLiC



Load code as shared lib, much faster:

```
.x mymacro.C+(42)
```

Uses the system's compiler, takes seconds

Subsequent `.x mymacro.C+(42)` check for changes, only rebuild if needed

Exactly as fast as e.g. Makefile based stand-alone binary!

CINT knows types, functions in the file, e.g. call

```
mymacro(43)
```





# Compiled versus Interpreted



 Why compile?

 Faster execution, CINT has limitations, validate code.

 Why interpret?

 Faster Edit → Run → Check result → Edit cycles ("rapid prototyping").

Scripting is sometimes just easier.

 Are Makefiles dead?

 Yes! ACLiC is even platform independent!

# A Little C[++]



Hopefully many of you know – but some don't.

- Object, constructor, assignment
- Pointer, reference
- Scope, destructor
- Stack vs. heap
- Inheritance, virtual functions

If you use C++ you *have* to understand these concepts!

# Objects, Constructors, =



Look at this code:

```
TNamed myObject("name", "title");  
TNamed mySecond;  
mySecond = myObject;  
cout << mySecond.GetName() << endl;
```

# Objects, Constructors, =



Look at this code:

```
TNamed myObject("name", "title");  
TNamed mySecond;  
mySecond = myObject;  
cout << mySecond.GetName() << endl;
```

Creating objects:

1. Constructor **TNamed::TNamed(const char\*, const char\*)**
2. Default constructor **TNamed::TNamed()**

# Objects, Constructors, =



Look at this code:

```
TNamed myObject("name", "title");  
TNamed mySecond;  
mySecond = myObject;  
cout << mySecond.GetName() << endl;
```



## 3. Assignment:

<i>mySecond</i>
TNamed: fName "" fTitle ""

<i>myObject</i>
TNamed: fName "name" fTitle "title"

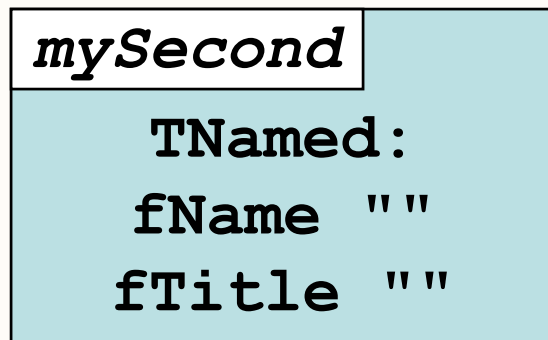
# Objects, Constructors, =



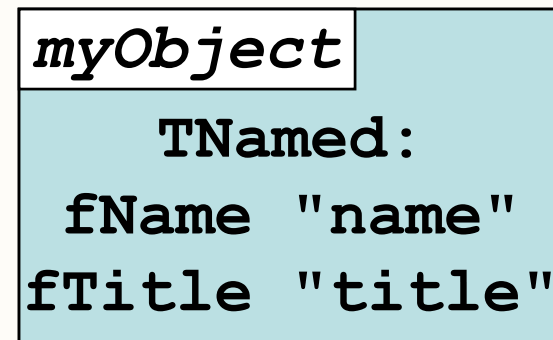
Look at this code:

```
TNamed myObject("name", "title");  
TNamed mySecond;  
mySecond = myObject;  
cout << mySecond.GetName() << endl;
```

## 3. Assignment: creating a twin



=



# Objects, Constructors, =



Look at this code:

```
TNamed myObject("name", "title");  
TNamed mySecond;  
mySecond = myObject;  
cout << mySecond.GetName() << endl;
```

## 4. New content

<i>mySecond</i>
TNamed: fName "name" fTitle "title"

output: "name"

# Pointer, Reference



Modified code:

```
TNamed myObject("name", "title");  
TNamed* pMySecond = 0;  
pMySecond = &myObject;  
cout << pMySecond->GetName() << endl;
```

Pointer declared with "\*", initialize to 0



# Pointer, Reference



Modified code:

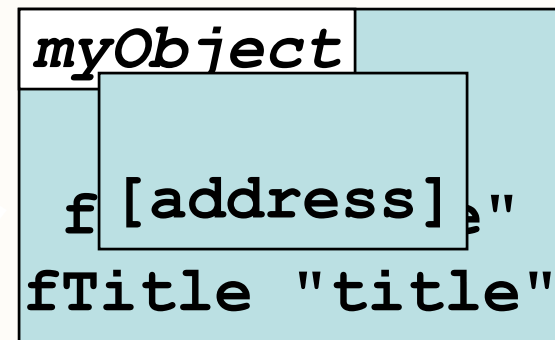
```
TNamed myObject("name", "title");  
TNamed* pMySecond = 0;  
pMySecond = &myObject;  
cout << pMySecond->GetName() << endl;
```

Assignment: point to myObject; no copy

*pMySecond*

=

&



# Pointer, Reference

Modified code:

```
TNamed myObject("name", "title");  
TNamed* pMySecond = 0;  
pMySecond = &myObject;  
cout << pMySecond->GetName() << endl;
```

Assignment: "&" creates reference:

<i>pMySecond</i>
[address]

=

&

<i>myObject</i>
TNamed: fName "name" fTitle "title"

# Pointer, Reference



Modified code:

```
TNamed myObject("name", "title");  
TNamed* pMySecond = 0;  
pMySecond = &myObject;  
cout << pMySecond->GetName() << endl;
```



Access members of value pointed to by "->"

# Pointer, Reference



Modified code:

```
TNamed myObject("name", "title");  
TNamed* pMySecond = 0;  
pMySecond = &myObject;  
cout << (*pMySecond).GetName() << endl;
```

Or dereference pointer by "\*" and then access like object with "."

Two blue arrows originate from the text below. One arrow points from the asterisk in "dereference pointer by '\*'" to the asterisk in "(\*pMySecond)" in the code. The other arrow points from the period in "access like object with '.'" to the period in ".GetName()" in the code.



# Pointer, Reference

Changes propagated:

```
TNamed myObject("name", "title");  
TNamed* pMySecond = 0;  
pMySecond = &myObject;  
pMySecond->SetName("newname");  
cout << myObject.GetName() << endl;
```

Pointer forwards to object

Name of object changed – prints "newname"!

# Object vs. Pointer



Compare object:

```
TNamed myObject("name", "title");  
TNamed mySecond = myObject;  
cout << mySecond.GetName() << endl;
```

to pointer:

```
TNamed myObject("name", "title");  
TNamed* pMySecond = &myObject;  
cout << pMySecond->GetName() << endl;
```

# Object vs. Pointer: Parameters



Calling functions: object parameter obj gets copied for function call!

```
void funcO(TNamed obj);  
TNamed myObject;  
funcO(myObject);
```

Pointer parameter: only address passed, no copy

```
void funcP(TNamed* ptr);  
TNamed myObject;  
funcP(&myObject);
```

# Object vs. Pointer: Parameters



Functions changing parameter: funcO can only access copy!  
**caller** not changed!

```
void funcO(TNamed obj) {  
    obj.SetName("nope");  
}  
  
funcO(caller);
```

Using pointers (or references) funcP can change  
**caller**

```
void funcP(TNamed* ptr) {  
    ptr->SetName("yes");  
}  
  
funcP(&caller);
```





# Scope

Scope: range of accessibility and C++ "life".

Birth: constructor, death: destructor

```
{ // birth: TNamed() called
  TNamed n;
} // death: ~TNamed() called
```

Variables are valid / accessible only in scopes:

```
int a = 42;
{ int a = 0; }
cout << a << endl;
```

# Scope



Functions are scopes:

```
void func() { TNamed obj; }  
  
func();  
cout << obj << end; // obj UNKNOWN!
```

must not return  
pointers to  
local variables!

```
TNamed* func() {  
    TNamed obj;  
    return &obj; // BAD!  
}
```

# Stack vs. Heap



So far only stack:

```
TNamed myObj ("n", "t");
```

Fast, but often < 10MB. Only survive in scope.

Heap: slower, GBs (RAM + swap), creation and destruction managed by user:

```
TNamed* pMyObj = new TNamed ("n", "t");  
delete pMyObj; // or memory leak!
```

# Stack vs. Heap: Functions



Can return heap objects without copying:

```
TNamed* CreateNamed() {  
    // user must delete returned obj!  
    TNamed* ptr = new TNamed("n", "t");  
    return ptr; }
```

ptr gone – but TNamed object still on the heap,  
address returned!

```
TNamed* pMyObj = CreateNamed();  
cout << pMyObj->GetName() << endl;  
delete pMyObj; // or memory leak!
```

# Inheritance



Classes "of same kind" can re-use functionality

E.g. TPlate, TBowl both dishes:

```
class TPlate: public TDish {...};  
class TBowl: public TDish {...};
```

Can implement common functions in TDish:

```
class TDish {  
    public:  
        void Wash();  
};
```

# Inheritance: The Base



Use TPlate, TBowl as dishes:

assign pointer of derived to pointer of base  
"every plate is a dish"

```
TDish *a = new TPlate();  
TDish *b = new TBowl();
```

But not every dish is a plate, i.e. the inverse  
doesn't work. And a bowl is totally not a plate!

```
TPlate* p = new TDish(); // NO!  
TPlate* q = new TBowl(); // NO!
```

# Virtual Functions



Often derived classes behave differently:

```
class TDish { ...
    virtual bool ForSoup() const;
};
class TPlate: public TDish { ...
    bool ForSoup() const {return false;}
};
class TBowl: public TDish { ...
    bool ForSoup() const {return true;}
};
```



# Pure Virtual Functions

But TDish cannot know! Mark as "not implemented"

```
class TDish { ...  
    virtual bool ForSoup() const = 0;  
};
```

Only for virtual functions.

Cannot create object of TDish anymore (one function is missing!)



# Calling Virtual Functions



Call to virtual functions evaluated at runtime:

```
void FillWithSoup(TDish* dish) {  
    if (dish->ForSoup())  
        dish->SetFull();  
}
```

Works for any type as expected:

```
TDish* a = new TPlate();  
TDish* b = new TBowl();  
FillWithSoup(a); // will not be full  
FillWithSoup(b); // is now full
```

# Virtual vs. Non-Virtual



So what happens if non-virtual?

```
class TDish { ...
    bool ForSoup() const {return false;}
};
```

Will now always call TDish::ForSoup(), i.e. false

```
void FillWithSoup(TDish* dish) {
    if (dish->ForSoup())
        dish->SetFull();
}
```

# Congrats!



*Fine print:*

*\* comes with lifelong  
subscription,*

*this diploma is worth nothing  
if not exercised regularly*

# Summary



We know:

- why and how to start ROOT
- C++ basics
- that you run your code with ".x"
- can call functions in libraries
- can (mis-) use ROOT as a pocket calculator!

Lots for you to discover during next two lectures  
and especially the exercises!



# Saving Data

Streaming, Reflection, TFile,  
Schema Evolution

# Saving Objects



Cannot do in C++:

```
TNamed* o; TNamed* p;  
o = new TNamed("name", "title");  
std::write("file.bin", "obj1", o);  
p = std::read("file.bin", "obj1");  
p->GetName();
```

E.g. LHC experiments use C++ to manage data  
Need to write C++ objects and read them back  
std::cout not an option: 15PetaBytes / year of  
processed data (i.e. data that will be read)

# Saving Objects – Saving Types



What's needed?

```
TNamed* o;  
o = new TNamed("name", "title");  
std::write("file.bin", "obj1", o);
```

Store *data members* of TNamed; need to know:

- 1) type of object
- 2) data members for the type
- 3) where data members are in memory
- 4) read their values from memory, write to disk

# Serialization



Store *data members* of TNamed: **serialization**

- 1) type of object: **runtime-type-information RTTI**
- 2) data members for the type: **reflection**
- 3) where data members are in memory:  
**introspection**
- 4) read their values from memory, write to disk:  
**raw I/O**

Complex task, and C++ is not your friend.



# Reflection



Need type description (aka *reflection*)

1. types, sizes, members

TMyClass is a class.

```
class TMyClass {  
    float fFloat;  
    Long64_t fLong;  
};
```

Members:

- "fFloat", type float, size 4 bytes
- "fLong", type Long64\_t, size 8 bytes

# Platform Data Types



Fundamental data types (int, long,...):  
size is platform dependent

Store "long" on 64bit platform, writing 8 bytes:

00, 00, 00, 00, 00, 00, 00, 42

Read on 32bit platform, "long" only 4 bytes:

00, 00, 00, 00

Data loss, data corruption!

# ROOT Basic Data Types



Solution: ROOT typedefs

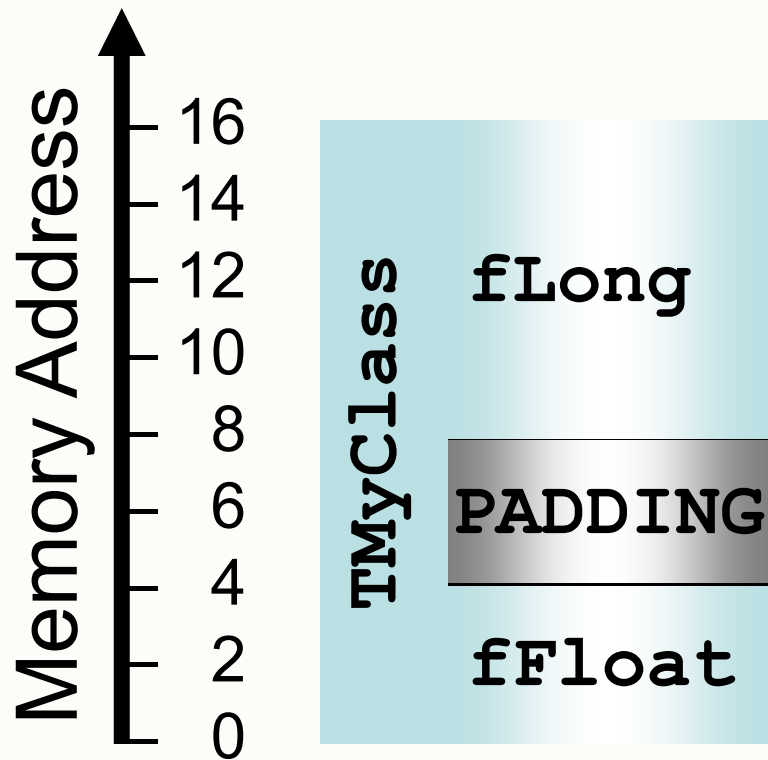
<b>Signed</b>	<b>Unsigned</b>	<b>sizeof [bytes]</b>
<code>Char_t</code>	<code>UChar_t</code>	1
<code>Short_t</code>	<code>UShort_t</code>	2
<code>Int_t</code>	<code>UInt_t</code>	4
<code>Long64_t</code>	<code>ULong64_t</code>	8
<code>Double32_t</code>		float on disk, double in RAM



# Reflection

Need type description (platform dependent)

1. types, sizes, members
2. offsets in memory



```
class TMyClass {  
    float fFloat;  
    Long64_t fLong;  
};
```

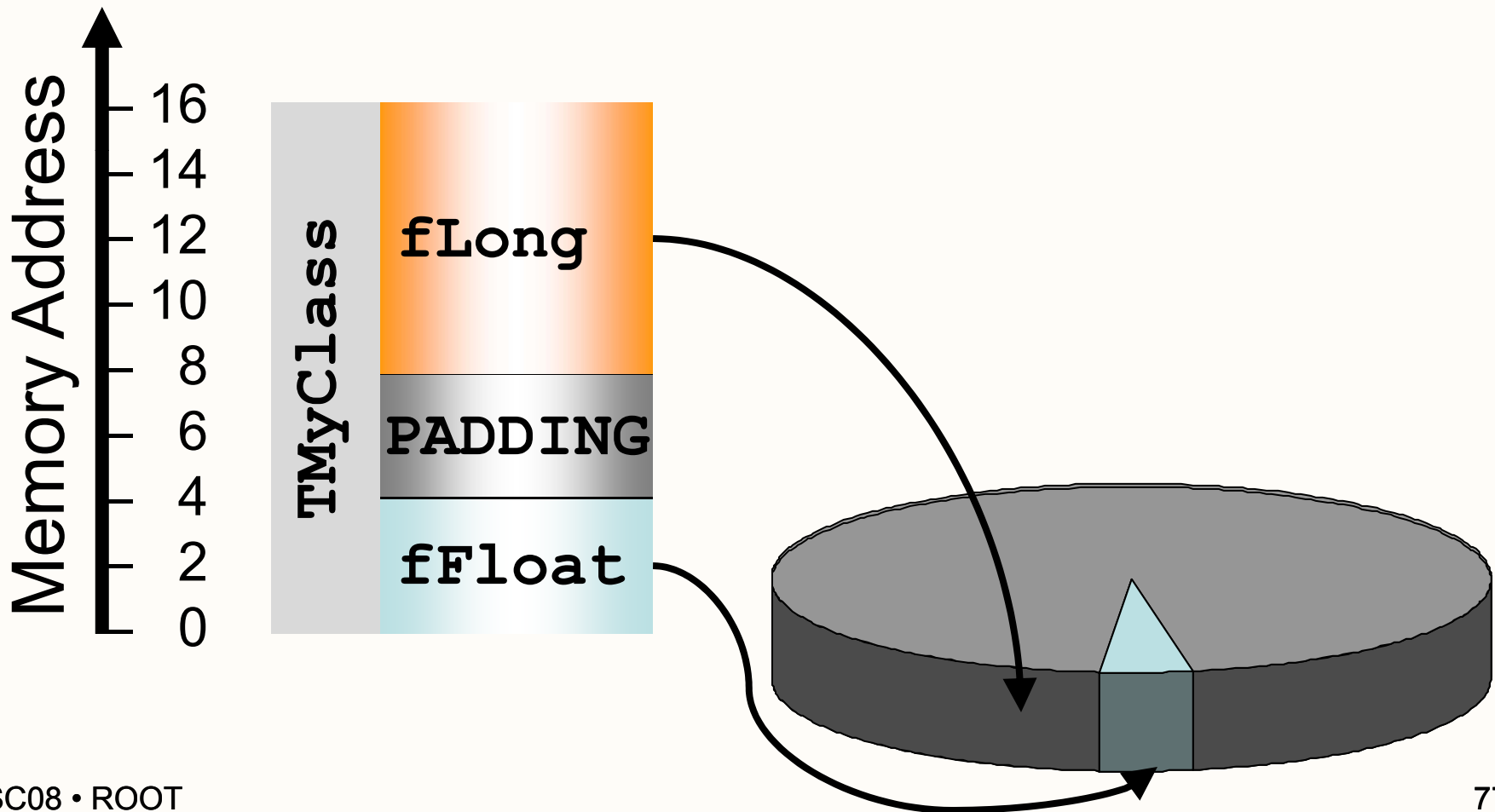
"fFloat" is at offset 0

"fLong" is at offset 8

# I/O Using Reflection



members → memory → disk



# C++ Is Not Java



Lesson: need reflection!

Where from?

Java: get data members with

```
Class.forName("MyClass").getFields()
```

C++: get data members with  
– oops. Not part of C++.

**BE CAREFUL**

**THIS LANGUAGE  
HAS NO BRAIN  
USE YOUR OWN**

# ROOT And Reflection



Simply use ACLiC:

```
.L MyCode.cxx+
```

Creates library with reflection data ("dictionary")  
of all types in MyCode.cxx!

Dictionary needed for interpreter, too  
ROOT has dictionary for all its types

# Back To Saving Objects: TFile



ROOT stores objects in TFiles:

```
TFile* f = new TFile("file.root", "NEW");
```

TFile behaves like file system:

```
f->mkdir("dir");
```

TFile has a current directory:

```
f->cd("dir");
```

TFile compresses data ("zip"):

```
f->GetCompressionFactor()  
2.61442160606384277e00
```



# Saving Objects, Really



Given a TFile:

```
TFile* f = new TFile("file.root", "RECREATE");
```

Write an object deriving from TObject:

```
object->Write("optionalName")
```

"optionalName" or TObject::GetName()

Write any object (with dictionary):

```
f->WriteObject(object, "name");
```



# "Where Is My Histogram?"

TFile owns histograms, graphs, trees  
(due to historical reasons):

```
TFile* f = new TFile("myfile.root");  
TH1F* h = new TH1F("h", "h", 10, 0., 1.);  
h->Write();  
TCanvas* c = new TCanvas();  
c->Write();  
delete f;
```

h automatically deleted: owned by file.

c still there.

→ *names unique!*

TFile acts like a scope for hists, graphs, trees!

# Risks With I/O



Physicists can loop a lot:

*For each particle collision*

*For each particle created*

*For each detector module*

*Do something.*

Physicists can loose a lot:

*Run for hours...*

*Crash.*

*Everything lost.*



# Name Cycles

Create snapshots regularly:

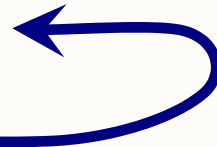
MyObject;1

MyObject;2

...

MyObject; 5427

MyObject



Write() does not replace but append!  
but see documentation TObject::Write()



# The "I" Of I/O

Reading is simple:

```
TFile* f = new TFile("myfile.root");  
TH1F* h = 0;  
f->GetObject("h", h);  
h->Draw();  
delete f;
```

Remember:

TFile owns histograms!

file gone, histogram gone!

# Ownership And TFiles



Separate TFile and histograms:

```
TFile* f = new TFile("myfile.root");  
TH1F* h = 0;  
TH1::AddDirectory(kFALSE);  
f->GetObject("h", h);  
h->Draw();  
delete f;
```

... and h will stay around.

# Changing Class – The Problem



Things change:

```
class TMyClass {  
    float fFloat;  
    Long64_t fLong;  
};
```

# Changing Class – The Problem



Things change:

```
class TMyClass {  
    double fFloat;  
    Long64_t fLong;  
};
```

Inconsistent reflection data, mismatch in  
memory, on disk

Objects written with old version cannot be read

*Need to store reflection with data to detect!*

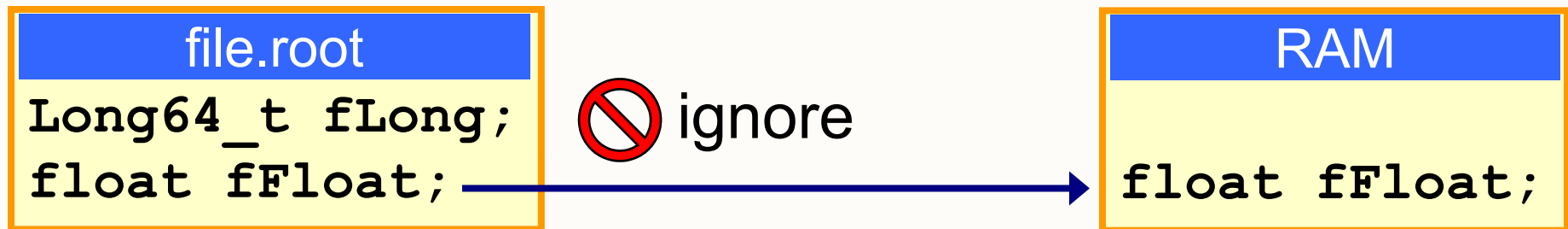


# Schema Evolution

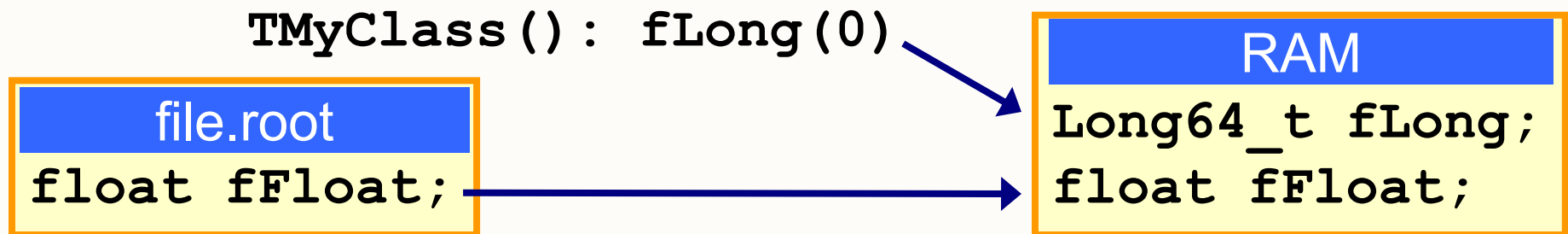


Simple rules to convert disk to memory layout

1. skip removed members



2. default-initialize added members



3. convert members where possible



# Class Version

ClassDef() macro makes I/O faster, needed when deriving from TObject

Can have multiple class versions in same file

Use version number to identify layout:

```
class TMyClass: public TObject {
public:
    TMyClass(): fLong(0), fFloat(0.) {}
    virtual ~TMyClass() {}
    ...
    ClassDef(TMyClass,1); // example class
};
```

# Reading Files



Files store reflection and data: need no library!

**With library**

event.fTemperature

Entries	1000
Mean	20.51
RMS	0.2901

*function*  
GetHistogram()  
*call*

**Without library**

event.fTemperature

Entries	1000
Mean	20.51
RMS	0.2901

# Powers of ROOT I/O



- Can even open  
`TFile("http://myserver.com/afile.root")`  
including read-what-you-need!
- Nice viewer for TFile: `new TBrowser`
- Combine contents of TFiles with  
`$ROOTSYS/bin/hadd`

# Summary



## Big picture:

- you know ROOT files – for petabytes of data
- you learned what schema evolution is
- you learned that reflection is key for I/O

## Small picture:

- you can write your own data to files
- you can read it back
- you can change the definition of your classes



# ROOT Collection Classes

# Collection Classes



ROOT collections polymorphic containers: hold pointers to **TObject**, so:

- Can only hold objects that inherit from **TObject**
- Return pointers to **TObject**, that have to be cast back to the correct subclass

```
void DrawHist(TObjArray *vect, int at)
{
    TH1F *hist = (TH1F*)vect->At(at);
    if (hist) hist->Draw();
}
```

# TObjArray



- Like `vector<TObject *>`, supports traditional array semantics:
  - Objects can be directly accessed via an index with the operator[]
  - Expands automatically when objects are added

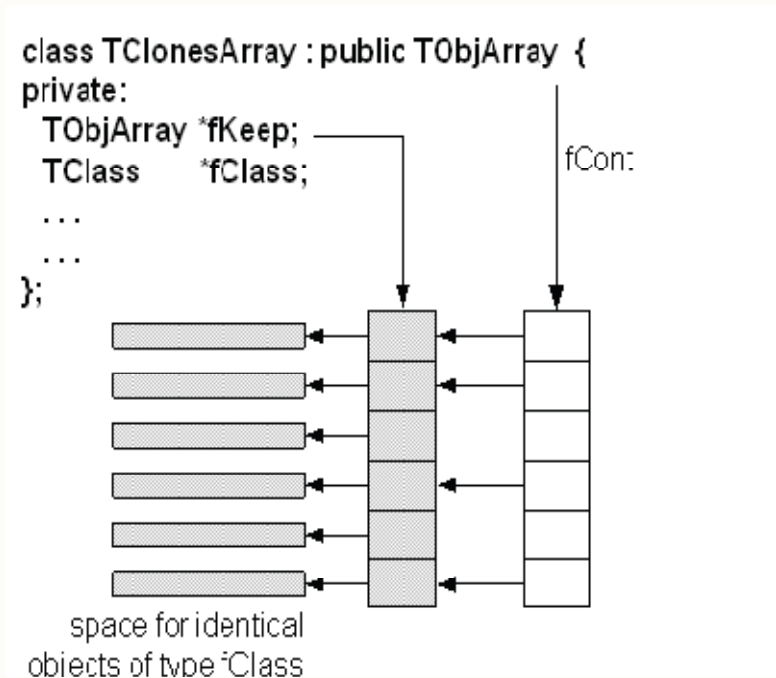


# TClonesArray

Array of objects of the same class ("clones")

Designed for repetitive data analysis tasks: same type of objects created and deleted many times.

No comparable class in STL!



*The internal data structure of a TClonesArray*



# Traditional Arrays

Very large number of new and delete calls in large loops like this (N(100000) x N(10000) times new/delete):

N(100000)

```
TObjArray a(10000);  
while (TEvent *ev = (TEvent *)next()) {  
    for (int i = 0; i < ev->Ntracks; ++i) {  
        a[i] = new TTrack(x,y,z,...);  
        ...  
    }  
    ...  
    a.Delete();  
}
```

N(10000)



You better use a **TClonesArray** which reduces the number of new/delete calls to only  $N(10000)$ :

```
TClonesArray a("TTrack", 10000);  
while (TEvent *ev = (TEvent *)next()) {  
    for (int i = 0; i < ev->Ntracks; ++i) {  
        new(a[i]) TTrack(x,y,z,...);  
        ...  
    }  
    ...  
    a.Delete();  
}
```

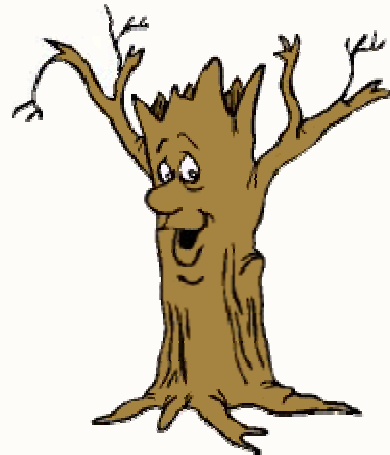
$N(100000)$

$N(10000)$

Pair of new / delete calls cost about  $4 \mu\text{s}$

Allocating / freeing memory  $NN(10^9)$  times costs about 1 hour!

# ROOT Trees



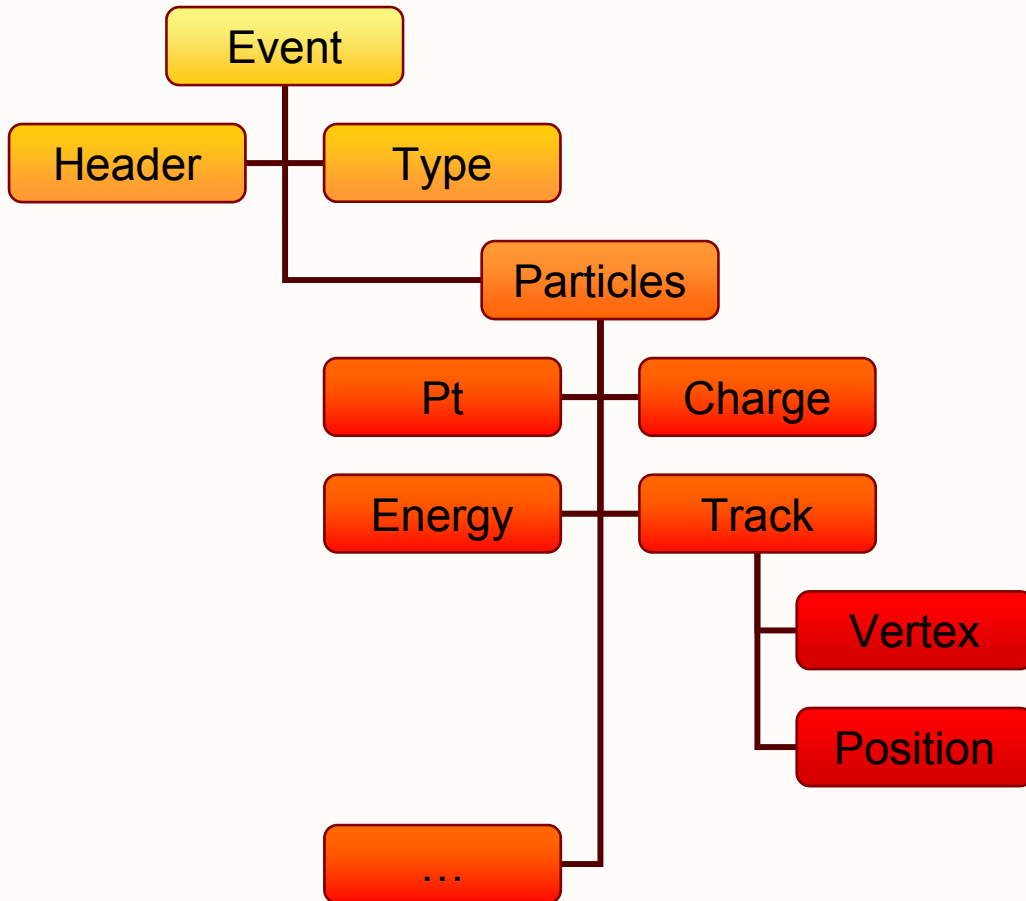


# Trees

From:  
Simple data types  
(e.g. Excel tables)

x	y	z
-1.10228	-1.79939	4.452822
1.867178	-0.59662	3.842313
-0.52418	1.868521	3.766139
-0.38061	0.969128	1.084074
0.552454	-0.21231	0.350281
-0.18495	1.187305	1.443902
0.205643	-0.77015	0.635417
1.079222	-0.32739	1.271904
-0.27492	-1.72143	3.038899
2.047779	-0.06268	4.197329
-0.45868	-1.44322	2.293266
0.304731	-0.88464	0.875442
-0.71234	-0.22239	0.556881
-0.27187	1.181767	1.470484
0.886202	-0.65411	1.213209
-2.03555	0.527648	4.421883
-1.45905	-0.464	2.344113
1.230661	-0.00565	1.514559
		3.562347

To:  
Complex data types  
(e.g. Database tables)



# Trees



- Databases have row wise access
  - Designed to store complete objects.
  - Data clustering is organized around objects and containers of objects.
  - Can only access the full object (e.g. full event)
- ROOT trees have column wise access
  - Direct access to any event, any branch or any leaf even in the case of variable length structures
  - Designed to access only a subset of the object attributes (e.g. only particles' energy)

# Why Trees ?



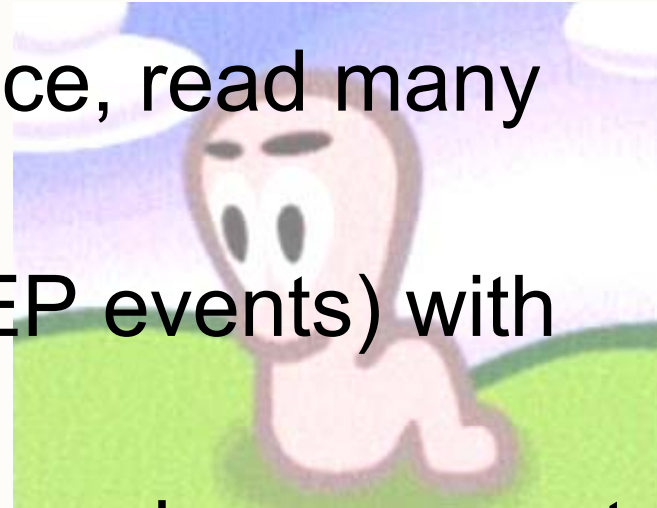
object.Write() convenient for simple objects like histograms, inappropriate for saving collections of events containing complex objects

- Reading a collection: read all elements (all events)
- With trees: only one element in memory, or even only a part of it (less I/O)

# Why Trees ?



- Extremely efficient write once, read many ("WORM")
- Designed to store  $>10^9$  (HEP events) with same data structure
- Trees allow fast direct and random access to any entry (sequential access is the best)
- Optimized for network access





# Building ROOT Trees



Overview of

- Trees
- Branches

5 steps to build a TTree

# Tree structure



ROOT Browser

File View Options

fTracks

All Folders

- root
- PROOF Sessions
- C:\home\bellenot\root\tutorials\tree
- ROOT Files
  - tree4.root
    - t4
      - event\_split
        - TObject
        - fEvtHdr
        - fTracks
        - fH
        - fTriggerBits
        - GetHistogram()

Contents of "/ROOT Files/tree4.root/t4/event\_split/fTracks"

Name	Title
fTracks.fBits	fBits[fTracks_]
fTracks.fBx	fBx[fTracks_]
fTracks.fBy	fBy[fTracks_]
fTracks.fCharge	fCharge[fTracks_]
fTracks.fMass2	fMass2[fTracks_]
fTracks.fMeanCharge	fMeanCharge[fTracks_]
fTracks.fNpoint	fNpoint[fTracks_]
fTracks.fNsp	fNsp[fTracks_]
fTracks.fPointValue	fPointValue[fTracks_]
fTracks.fPx	fPx[fTracks_]
fTracks.fPy	fPy[fTracks_]
fTracks.fPt	fPt[fTracks_]

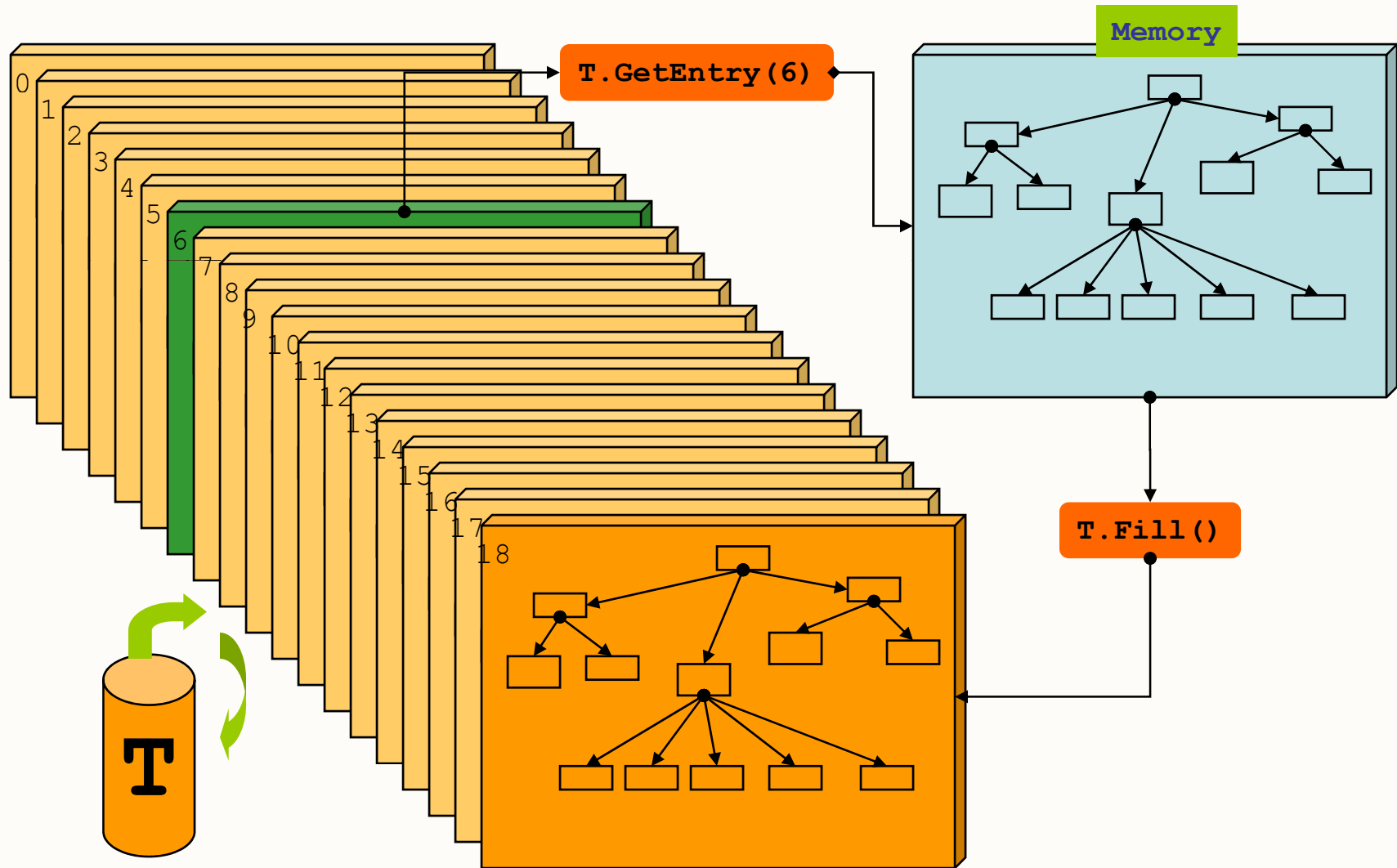
# Tree structure



- Branches: directories
- Leaves: data containers
- Can read a subset of all branches – speeds up considerably the data analysis processes
- Tree layout can be optimized for data analysis
- The class for a branch is called **TBranch**
- Variables on **TBranch** are called leaf (yes - **TLeaf**)
- Branches of the same **TTree** can be written to separate files

# Memory ↔ Tree

Each Node is a branch in the Tree



# Five Steps to Build a Tree



## Steps:

1. Create a TFile
2. Create a TTree
3. Add TBranch to the TTree
4. Fill the tree
5. Write the file

# Example macro

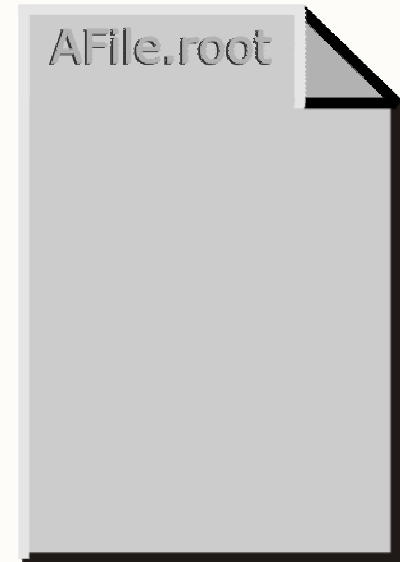


```
void WriteTree()
{
    Event *myEvent = new Event();
    TFile f("AFile.root");
    TTree *t = new TTree("myTree", "A Tree");
    t->Branch("EventBranch", myEvent);
    for (int e=0; e<100000; ++e) {
        myEvent->Generate(); // hypothetical
        t->Fill();
    }
    t->Write();
}
```

# Step 1: Create a TFile Object



Trees can be huge → need file  
for swapping filled entries



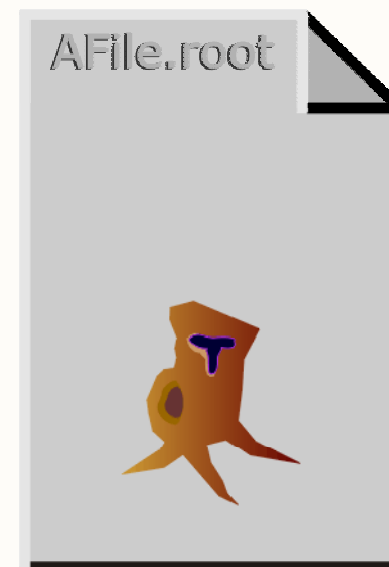
```
TFile *hfile = new TFile("AFile.root");
```

# Step 2: Create a TTree Object



The TTree constructor:

- Tree name (e.g. "myTree")
- Tree title



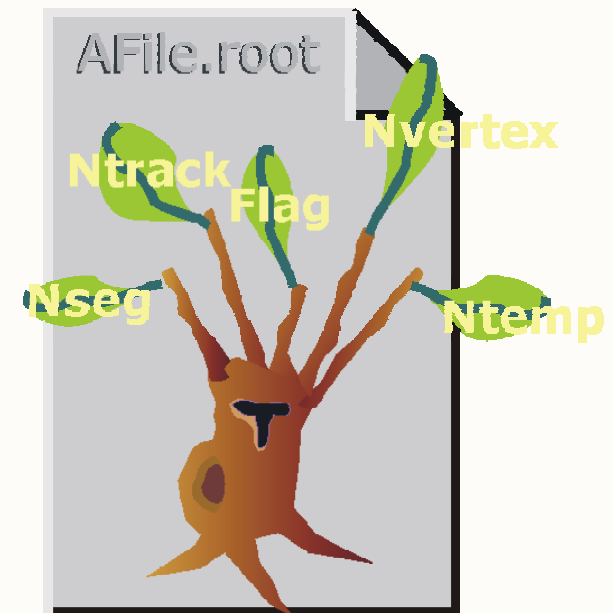
```
TTree *tree = new TTree("myTree", "A Tree");
```



# Step 3: Adding a Branch



- Branch name
- Pointer to the object

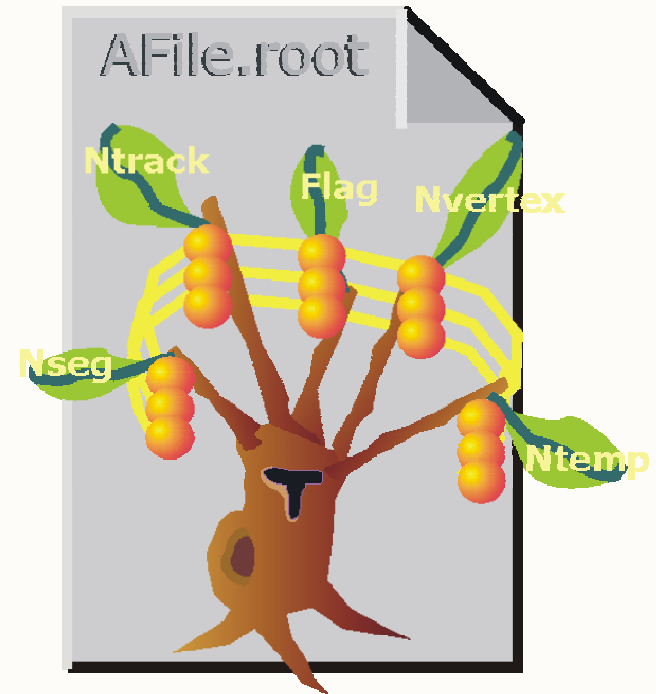


```
Event *myEvent = new Event();  
myTree->Branch("eBranch", myEvent);
```

# Step 4: Fill the Tree



- Create a for loop
- Assign values to the object contained in each branch
- TTree::Fill() creates a new entry in the tree: snapshot of values of branches' objects

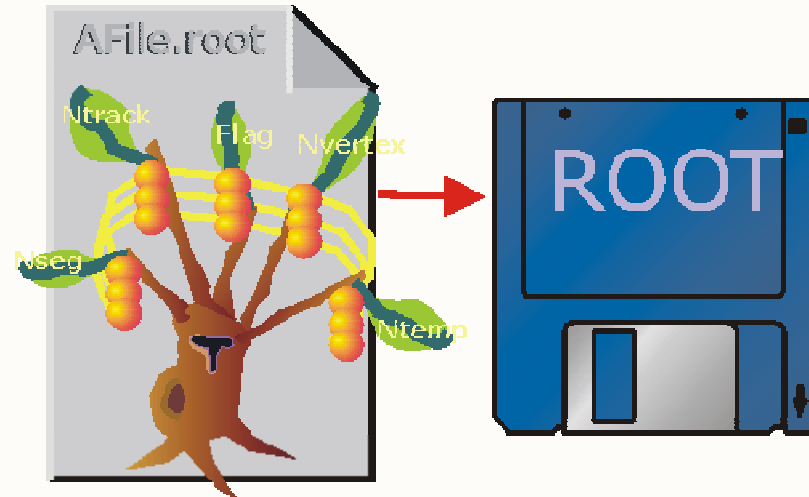


```
for (int e=0;e<100000;++e) {  
    myEvent->Generate(e); // fill event  
    myTree->Fill();      // fill the tree  
}
```

# Step 5: Write Tree To File



```
myTree->Write();
```



# Reading a TTree



- Looking at a tree
- How to read a tree
- Friends and chains

# Example macro



```
void ReadTree()  
{  
    Event *myEvent = 0;  
    TFile f("AFile.root");  
    TTree *myTree = (TTree*) f->Get("myTree");  
    myTree->SetBranchAddresses("EventBranch",  
                               myEvent);  
    for (int e=0; e<100000; ++e) {  
        myTree->GetEntry(e);  
        myEvent->Analyze();  
    }  
}
```



The pointer (myEvent) MUST be set to 0

# How to Read a TTree



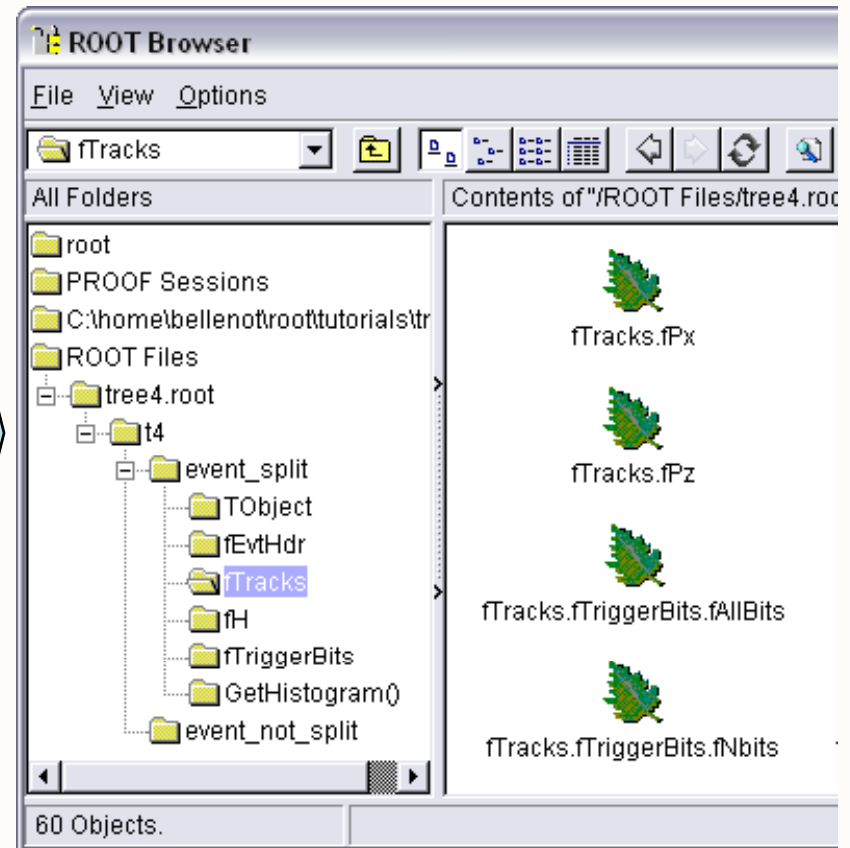
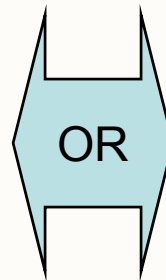
Example:

1. Open the Tfile

```
TFile f("AFile.root")
```

2. Get the TTree

```
TTree *myTree = 0;  
f.GetObject("myTree",  
myTree)
```



# How to Read a TTree



3. Create a variable pointing to the data

```
root [] Event *myEvent = 0;
```

4. Associate a branch with the variable:

```
root [] myTree->SetBranchAddress("eBranch", myEvent);
```

5. Read one entry in the TTree

```
root [] myTree->GetEntry(0)
```

```
root [] myEvent->GetTracks()->First()->Dump()
```

```
==> Dumping object at: 0x0763aad0, name=Track,  
    class=Track
```

```
fPx          0.651241      X component of the momentum  
fPy          1.02466      Y component of the momentum  
fPz          1.2141       Z component of the momentum  
[...]
```

# Branch Access Selection



- Use `TTree::SetBranchStatus()` to activate only the branches holding wanted variables.
- Speed up considerably the reading phase

```
TClonesArray* myMuons = 0;  
// disable all branches  
myTree->SetBranchStatus("*", 0);  
// re-enable the "muon" branches  
myTree->SetBranchStatus("muon*", 1);  
myTree->SetBranchAddress("muon", myMuons);  
// now read (access) only the "muon" branches  
myTree->GetEntry(0);
```



# Looking at the Tree



TTree::Print() shows the data layout

```
root [] TFile f("AFile.root")
root [] myTree->Print();
```

```
*****
*Tree      :myTree      : A ROOT tree *
*Entries   :          10 : Total =          867935 bytes File Size =          390138 *
*          :            : Tree compression factor =          2.72 *
*****
*Branch    :eBranch *
*Entries   :          10 : BranchElement (see below) *
*..... *
*Br       0 :fUniqueID : *
*Entries   :          10 : Total Size=          698 bytes One basket in memory *
*Baskets   :           0 : Basket Size=          64000 bytes Compression=          1.00 *
*..... *
...
...
```

# Looking at the Tree



TTree::Scan("leaf:leaf:....") shows the values

```
root [] myTree->Scan("fNseg:fNtrack"); > scan.txt
```

```
root [] myTree->Scan("fEvtHdr.fDate:fNtrack:fPx:fPy", "",  
"colsize=13 precision=3 col=13:7::15.10");
```

```
*****  
* Row * Instance * fEvtHdr.fDate * fNtrack *          fPx *          fPy *  
*****  
*  0 *         0 *      960312 *    594 *          2.07 *      1.459911346 *  
*  0 *         1 *      960312 *    594 *          0.903 *     -0.4093382061 *  
*  0 *         2 *      960312 *    594 *          0.696 *      0.3913401663 *  
*  0 *         3 *      960312 *    594 *         -0.638 *      1.244356871 *  
*  0 *         4 *      960312 *    594 *         -0.556 *     -0.7361358404 *  
*  0 *         5 *      960312 *    594 *         -1.57 *     -0.3049036264 *  
*  0 *         6 *      960312 *    594 *          0.0425 *     -1.006743073 *  
*  0 *         7 *      960312 *    594 *         -0.6 *     -1.895804524 *  
*****
```

# TTree Selection Syntax



```
MyTree->Scan ();
```

Prints the first 8 variables of the tree.

```
MyTree->Scan ("*");
```

Prints all the variables of the tree.

Select specific variables:

```
MyTree->Scan ("var1:var2:var3");
```

Prints the values of var1, var2 and var3.

A selection can be applied in the second argument:

```
MyTree->Scan ("var1:var2:var3", "var1>0");
```

Prints the values of var1, var2 and var3 for the entries where var1 is greater than 0

Use the same syntax for TTree::Draw()

# Looking at the Tree



`TTree::Show(entry_number)` shows the values for one entry

```
root [] myTree->Show(0);
=====> EVENT:0
eBranch          = NULL
fUniqueID        = 0
fBits            = 50331648
[...]
fNtrack          = 594
fNseg            = 5964
[...]
fEvtHdr.fRun     = 200
[...]
fTracks.fPx      = 2.066806, 0.903484, 0.695610, -0.637773, ...
fTracks.fPy      = 1.459911, -0.409338, 0.391340, 1.244357, ...
```

# TChain: the Forest



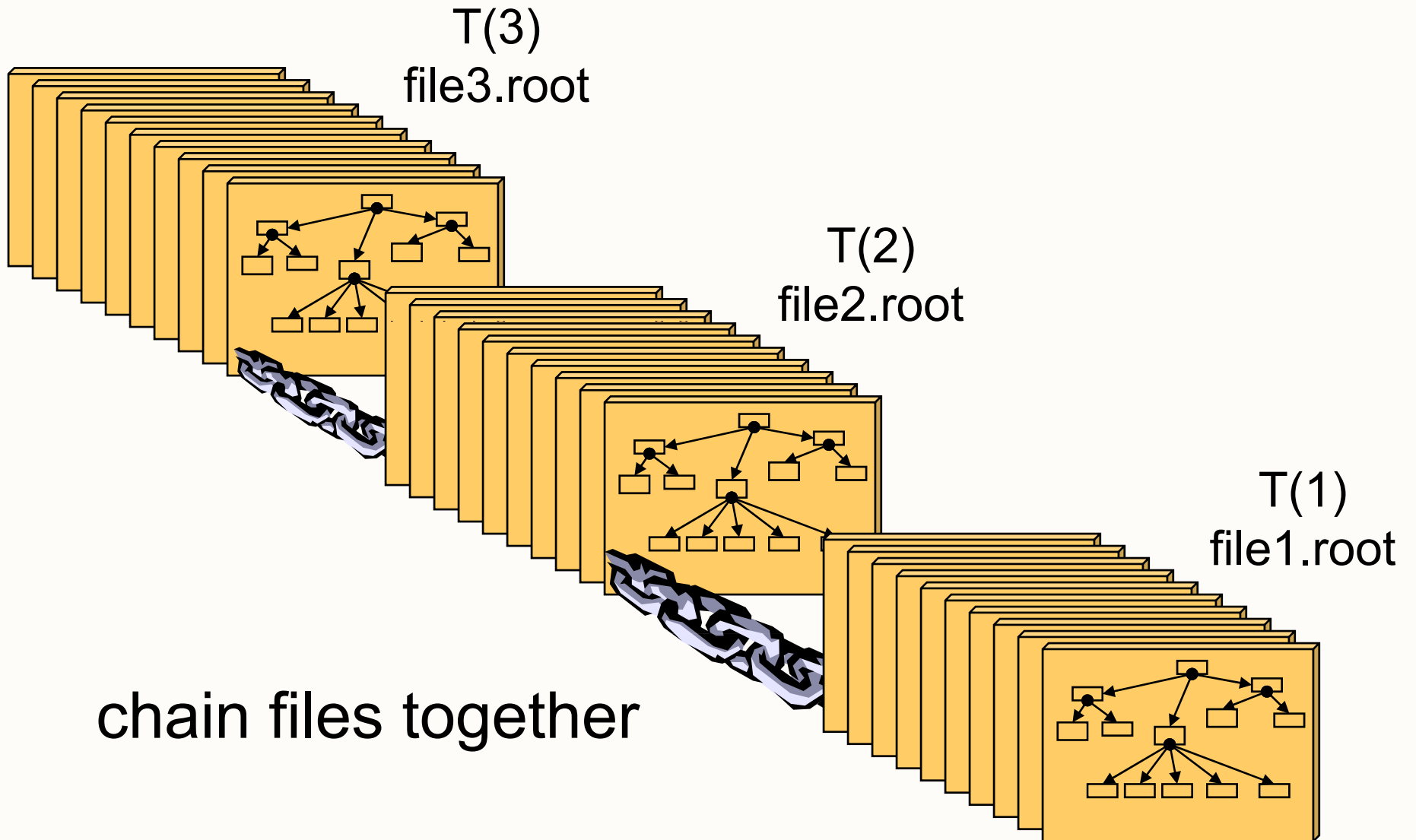
- Collection of TTrees: list of ROOT files containing the same tree
- Same semantics as TTree

As an example, assume we have three files called file1.root, file2.root, file3.root. Each contains tree called "T". Create a chain:

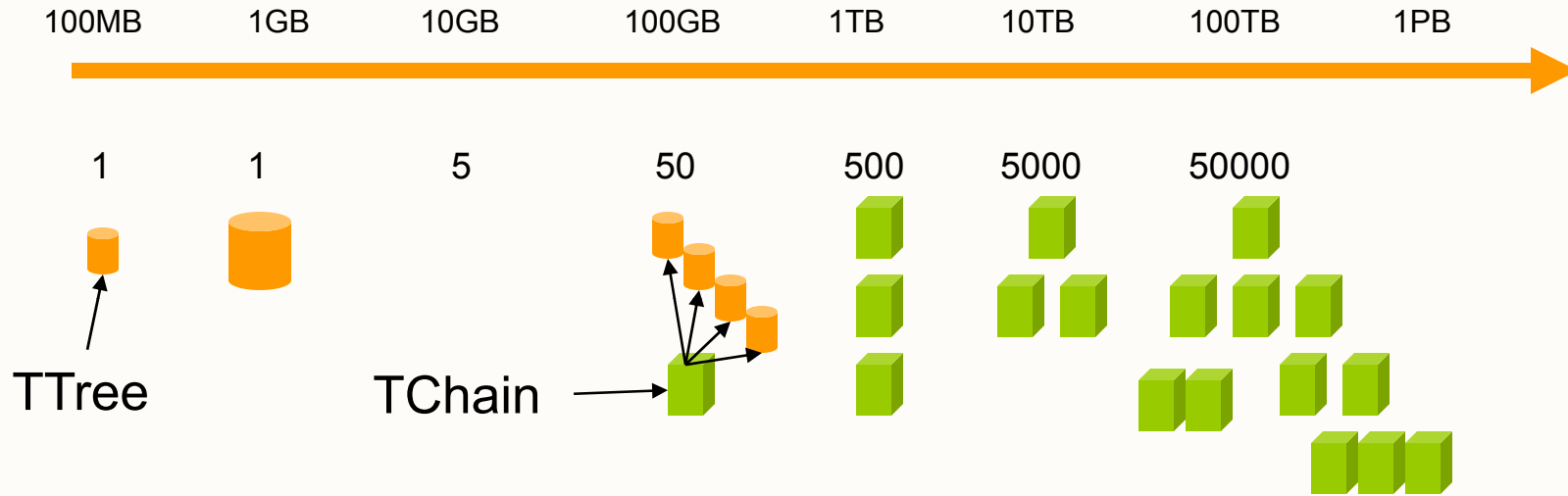
```
TChain chain("T"); // argument: tree name  
chain.Add("file1.root");  
chain.Add("file2.root");  
chain.Add("file3.root");
```

Now we can use the TChain like a TTree!

# TChain



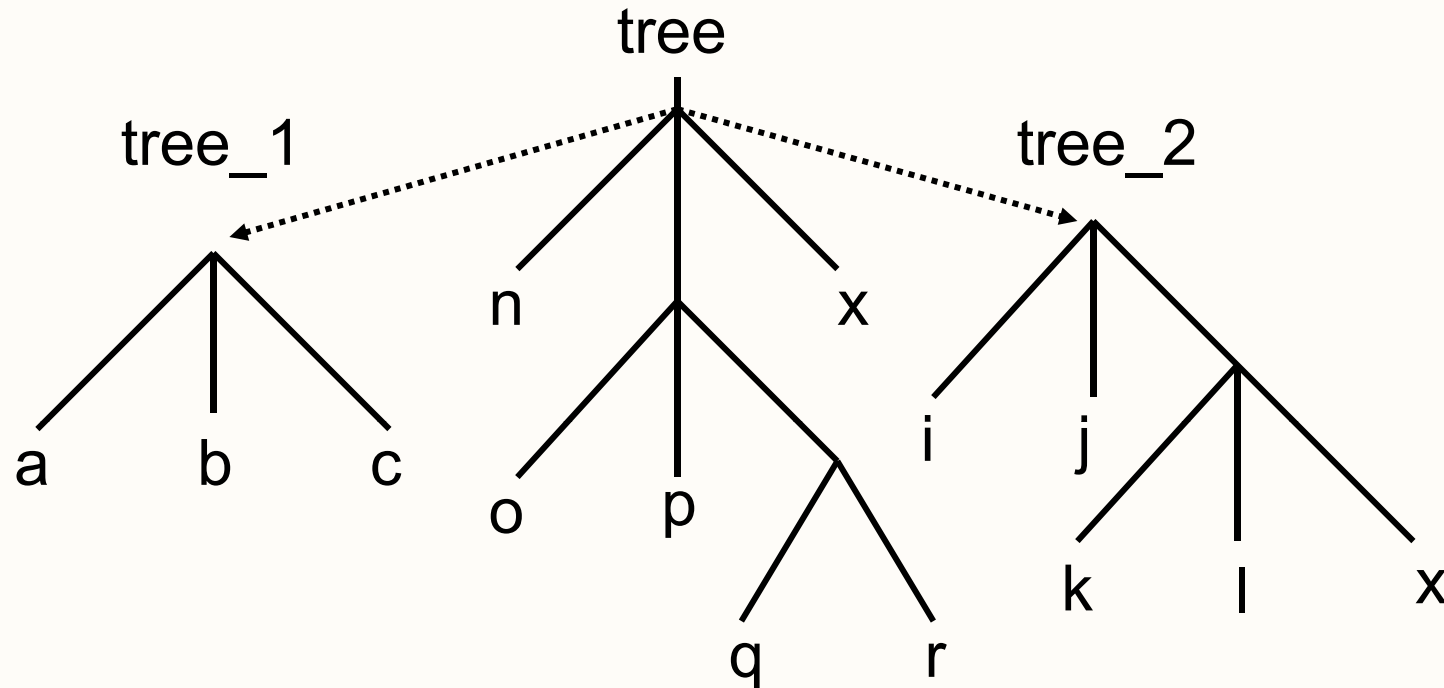
# Data Volume & Organisation



- A TFile typically contains 1 TTree
- A TChain is a collection of TTrees or/and TChains
- A TChain is typically the result of a query to a file catalog



# Tree Friends



```
TFile f1("tree1.root");  
tree.AddFriend("tree_1", "tree2.root");  
tree.AddFriend("tree_2", "tree3.root");  
tree.Draw("x:a", "k<c");  
tree.Draw("x:tree_2.x", "sqrt(p)<b");
```

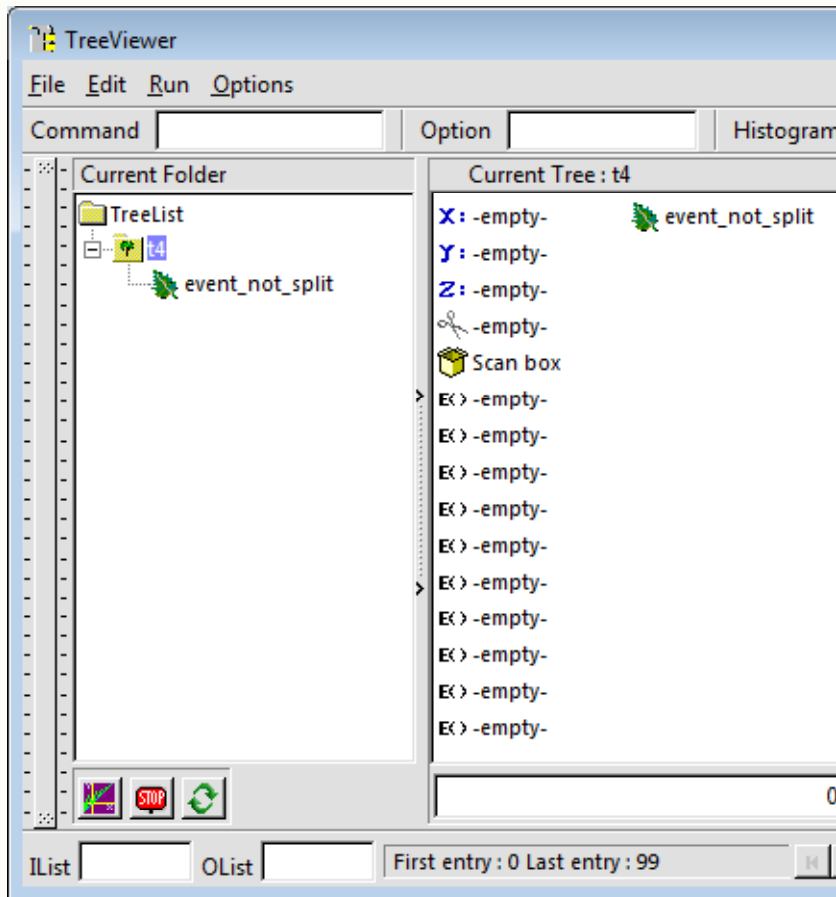


# Summary: Trees, basics

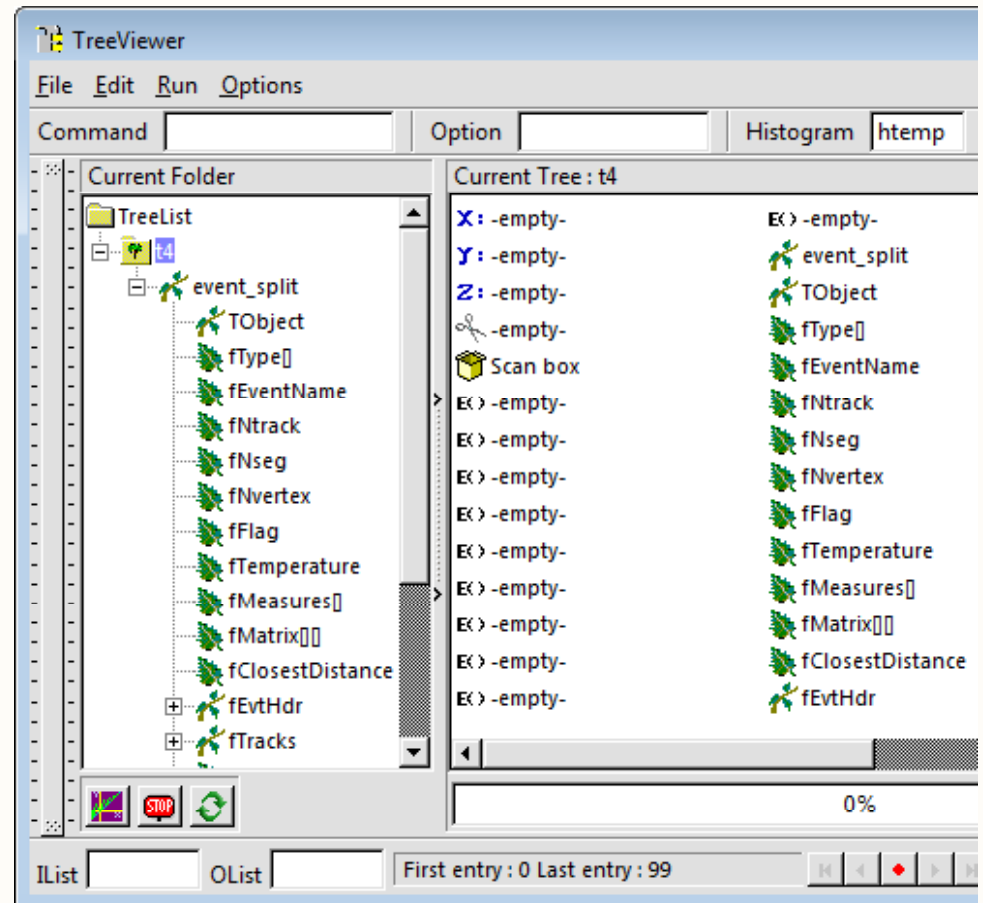


- TTree is one of the most powerful collections available for HEP
- Extremely efficient for huge number of data sets with identical layout
- Very easy to look at TTree - use TBrowser!
- Write once, read many (WORM) ideal for experiments' data
- Still: extensible, users can add their own tree as friend

# Splitting



Split level = 0



Split level = 99

# Splitting



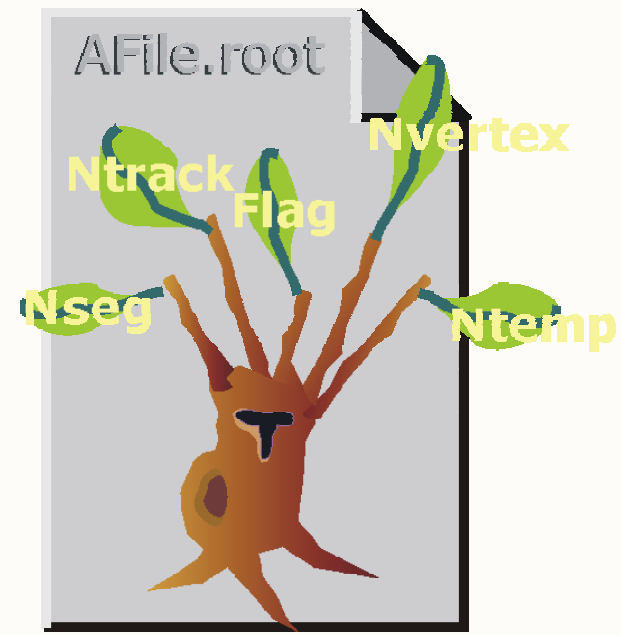
- Creates one branch per member – recursively
- Allows to browse objects that are stored in trees, even without their library
- Makes same members consecutive, e.g. for object with position in X, Y, Z, and energy E, all X are consecutive, then come Y, then Z, then E. A lot higher zip efficiency!
- Fine grained branches allow fine-grained I/O - read only members that are needed, instead of full object
- Supports STL containers, too!

# Splitting

Setting the split level (default = 99)



Split level = 0



Split level = 99

```
tree->Branch("EvBr", &event, 64000, 0);
```

# Performance Considerations



A split branch is:

- Faster to read - the type does not have to be read each time
- Slower to write due to the large number of branches

# Analyzing Trees

Selectors, Analysis, PROOF



# Recap



TTree efficient storage and access  
for huge amounts of structured data

Allows selective access of data

TTree knows its layout

Almost all HEP analyses based on TTree

# TTree Data Access



TSelector: generic "TTree based analysis"

Derive from it ("TMySelector")

ROOT invokes TSelector's functions,

Used e.g. by `tree->Process(TSelector*,...),`

PROOF

Functions called are virtual, thus TMySelector's functions called.





# TSelector

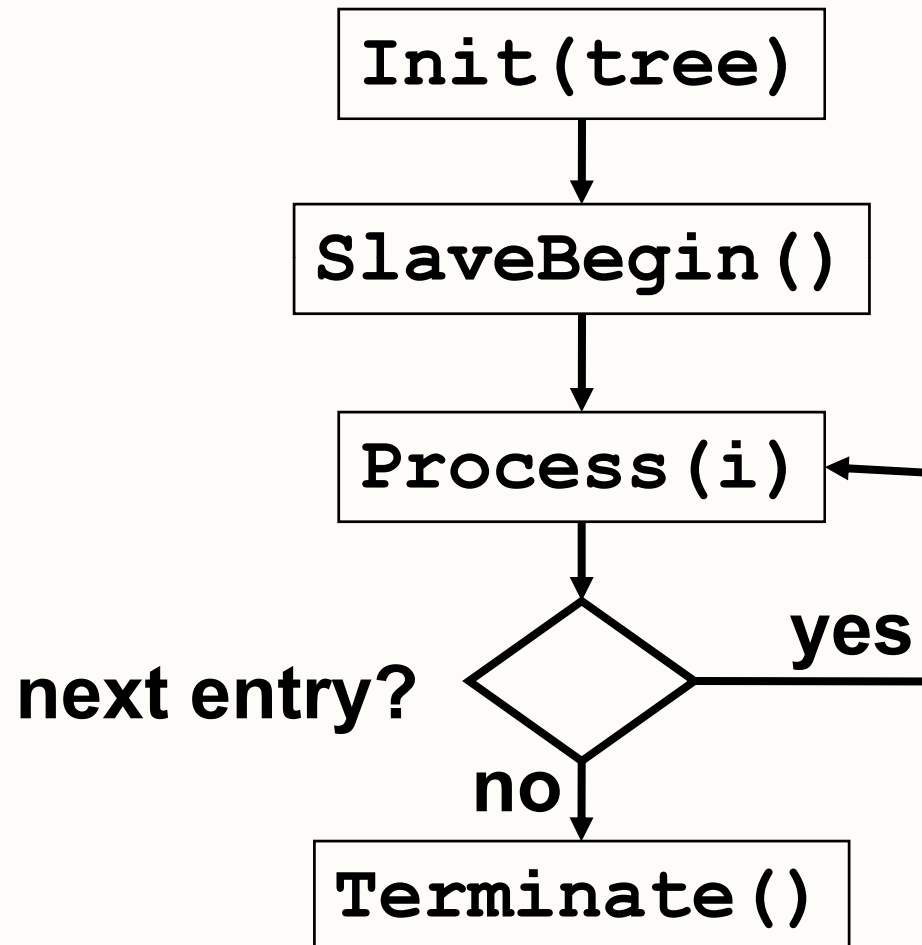
Steps of ROOT using a TSelector:

- 1. *setup*** `Init(TTree*)`  
called to inform selector about tree
- 2. *start*** `SlaveBegin()`  
called to create histograms
- 3. *run*** `Process(Long64_t)`  
called for each entry to load and analyze it
- 4. *end*** `Terminate()`  
called to fit histograms, write them to files,...

# TTree Data Access



E.g. `tree->Process ("MySelector.C+")`





# TSelector: Usage

- `Init(TTree* tree):`  
e.g. `TMySelector::fChain = tree.`  
Set branch addresses.
- `SlaveBegin():` create histograms
- `Process(Long64_t entry):`  
`fChain->GetTree()->GetEntry(entry);`  
fill histograms
- `Terminate():` fit; save histograms

# Analysis Example

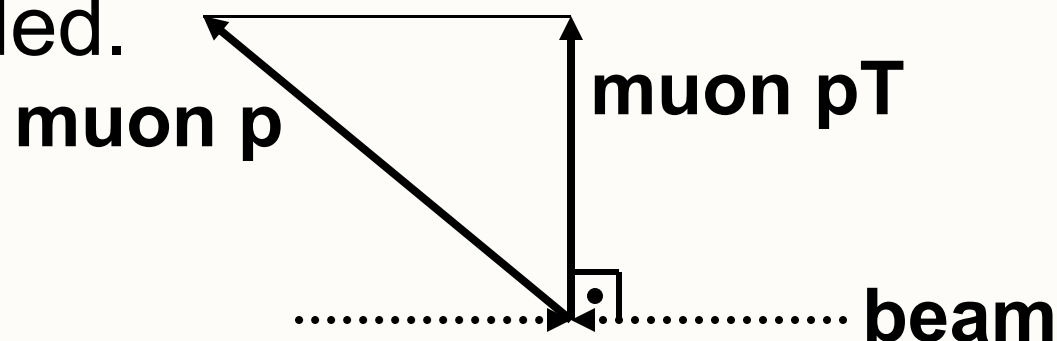


Determine trigger efficiencies from data, typical ingredient in analyses

Trigger selection before writing data: not all events available

Usually higher energy is taken, lower is ignored

Example 15GeV muon trigger: events with a muon  $> 15\text{GeV}$  transverse momentum ("pT") are recorded.



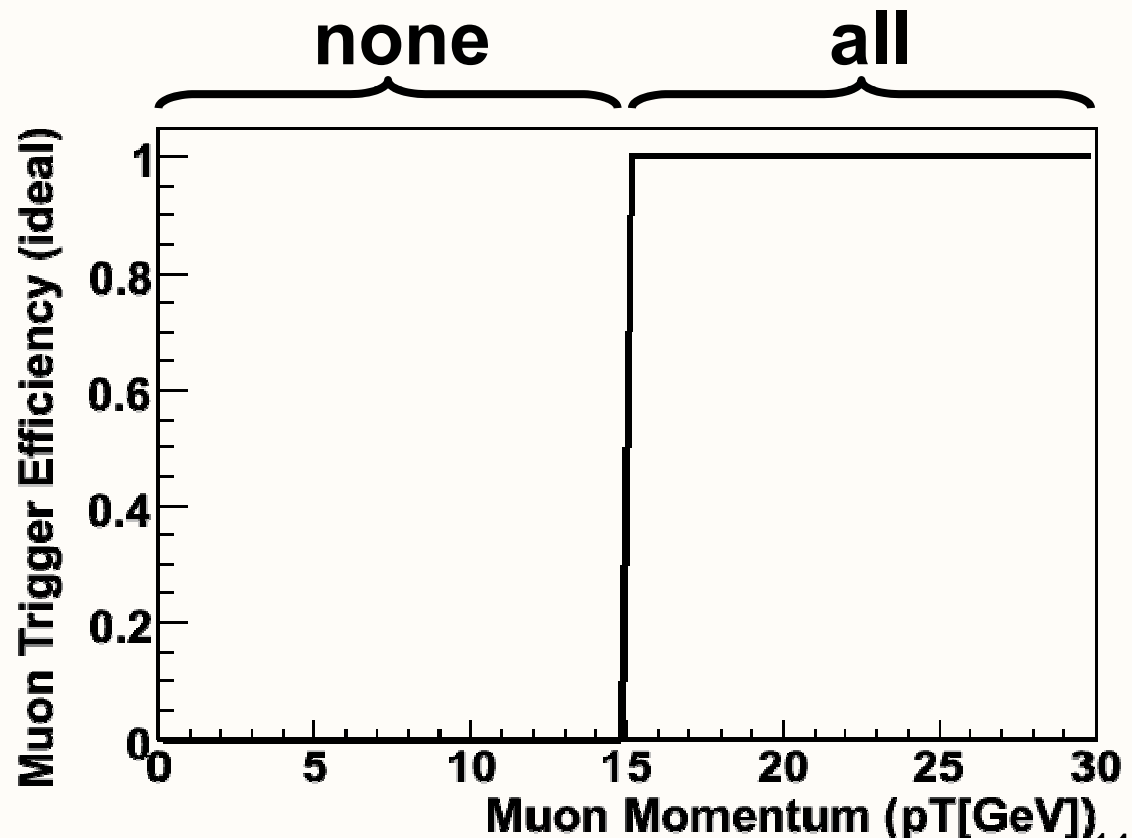
# Ideal Trigger



Efficiency: probability to record an event with a given (transverse) muon momentum

$$\text{eff} = \text{triggered}/\text{all}$$

Ideal 15GeV  
muon trigger



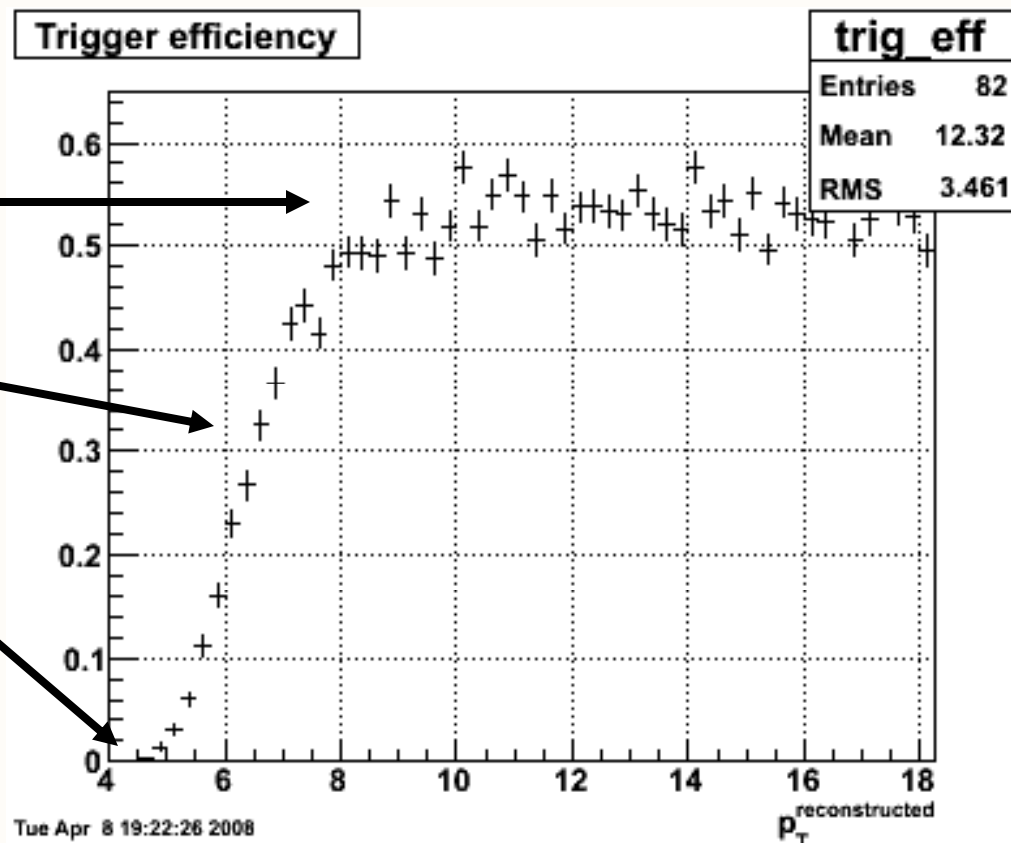
# Trigger Example



Example for a trigger from STAR @ BNL

Main properties:

- plateau
- turn-on
- minimum



# Trigger Efficiency From Data



Look at data triggered by 15GeV muon trigger:  
for each event's muon:

T=triggered, N=not triggered

{T} {NTN}{TT}{TNT}{NT}{TT}...

But this sample doesn't see {N}, {NN}, {NNN},...

$$\text{eff} = |T| / \text{all} = |T| / (|T| + |N|)$$

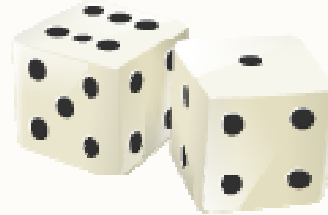
But |N| unknown! Cannot determine efficiency!

Instead: need muons that are independent of  
trigger ("unbiased")

# Dice And Tag / Probe



Think of two dice



Want probability for "6" ("6" trigger efficiency)

Have only triggered data, all results have  $>1$  "6"

$\{1,6\}, \{6,4\}, \{6,6\}, \{1,6\}, \dots$

Solution: one die has 6, the other is unbiased!

Result: N, N, T, N, ... will yield  $1/6$



# Muons And Tag / Probe



Solution: events with  $>1$  muon

For each muon:

if exists other muon causing trigger:  
this muon is unbiased!

Need trigger decision stored with data, as in:

"other muon caused the 15GeV muon trigger"

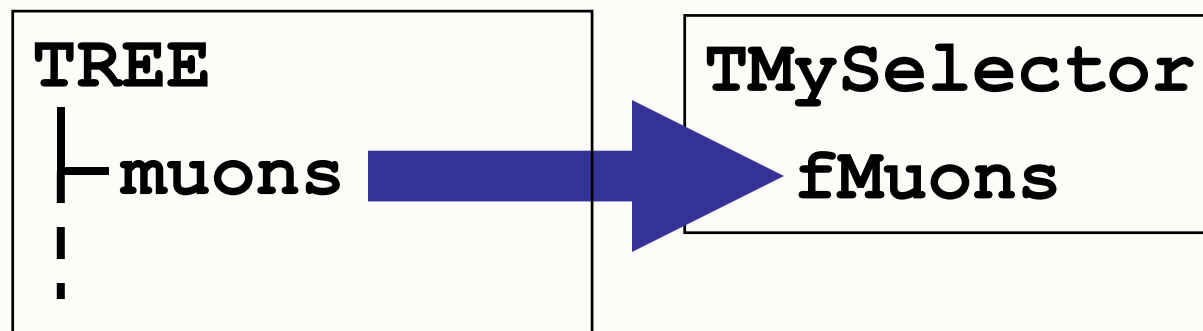


# Get Data From TTree

In TSelector::Init(tree) select branches and connect tree with our member fMuons:

```
TTree* t = fChain->GetTree();  
t->SetBranchStatus("*", 0); // all off  
t->SetBranchStatus("muons*", 1); // but muons  
t->SetBranchAddress("muons", fMuons);
```

TTree::GetEntry(i) will load data from branch muons into fMuons; can access data via fMuons





# TClonesArray

fMuons could be TClonesArray:

```
TClonesArray* fMuons; // array of TMuon
class TMuon: public TObject {
public:
    ...
    float Pt() const;
    bool Mu15() const; // triggered
};
```

Print pT of the i-th muon:

```
TMuon* muon = (TMuon*) fMuons->At(i);
cout << muon->Pt() << endl;
```



# Determine Efficiency

Take a random muon number  $i$  ("probe")

Check that another muon ("tag") has caused trigger, **then**:

```
++ all[probe->pT()]
```

```
if probe muon has triggered:
```

```
++fired[probe->pT()]
```

```
efficiency[pT] = fired[pT] / all[pT]
```

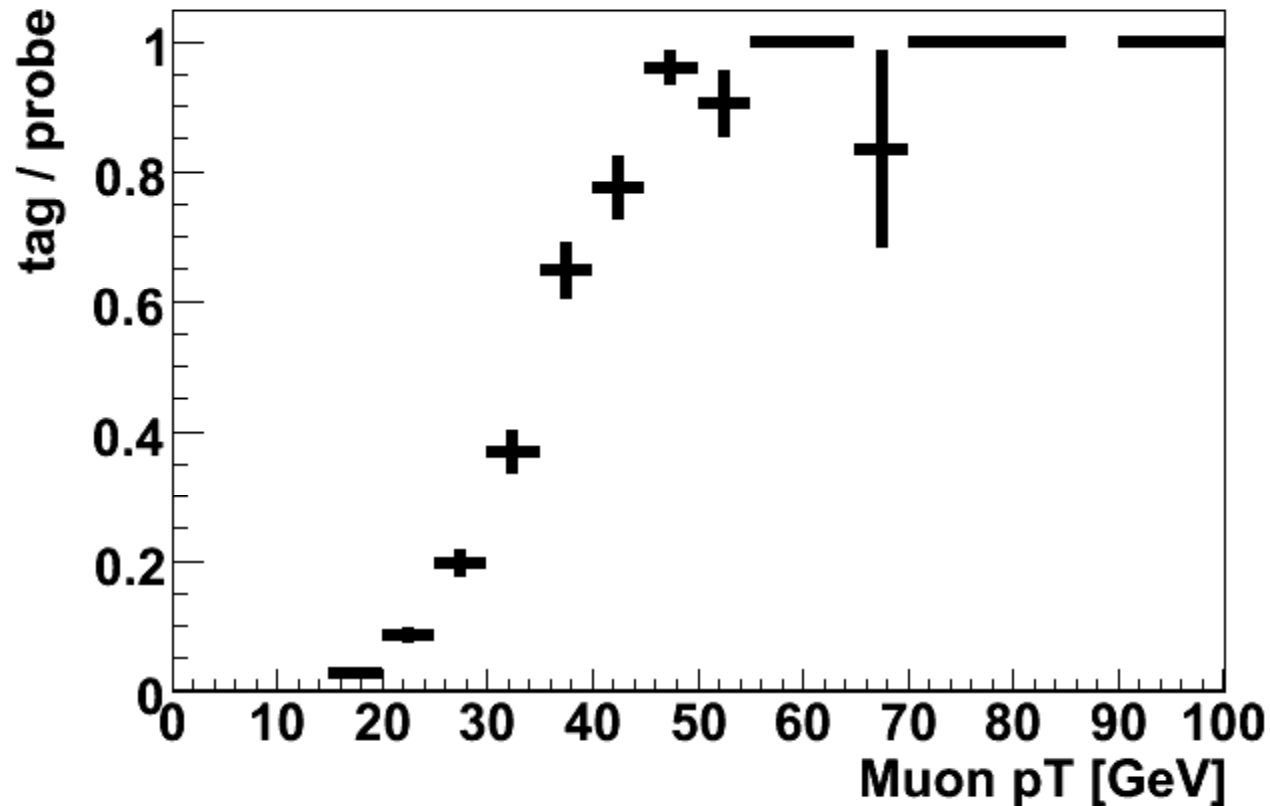
Counting in  $pT$ -bins – use histogram

Division: binomial errors, check Wikipedia ;-)

# Result



Dividing probes / tags yields sampled efficiency  
"Bumpy" because of low numbers of events

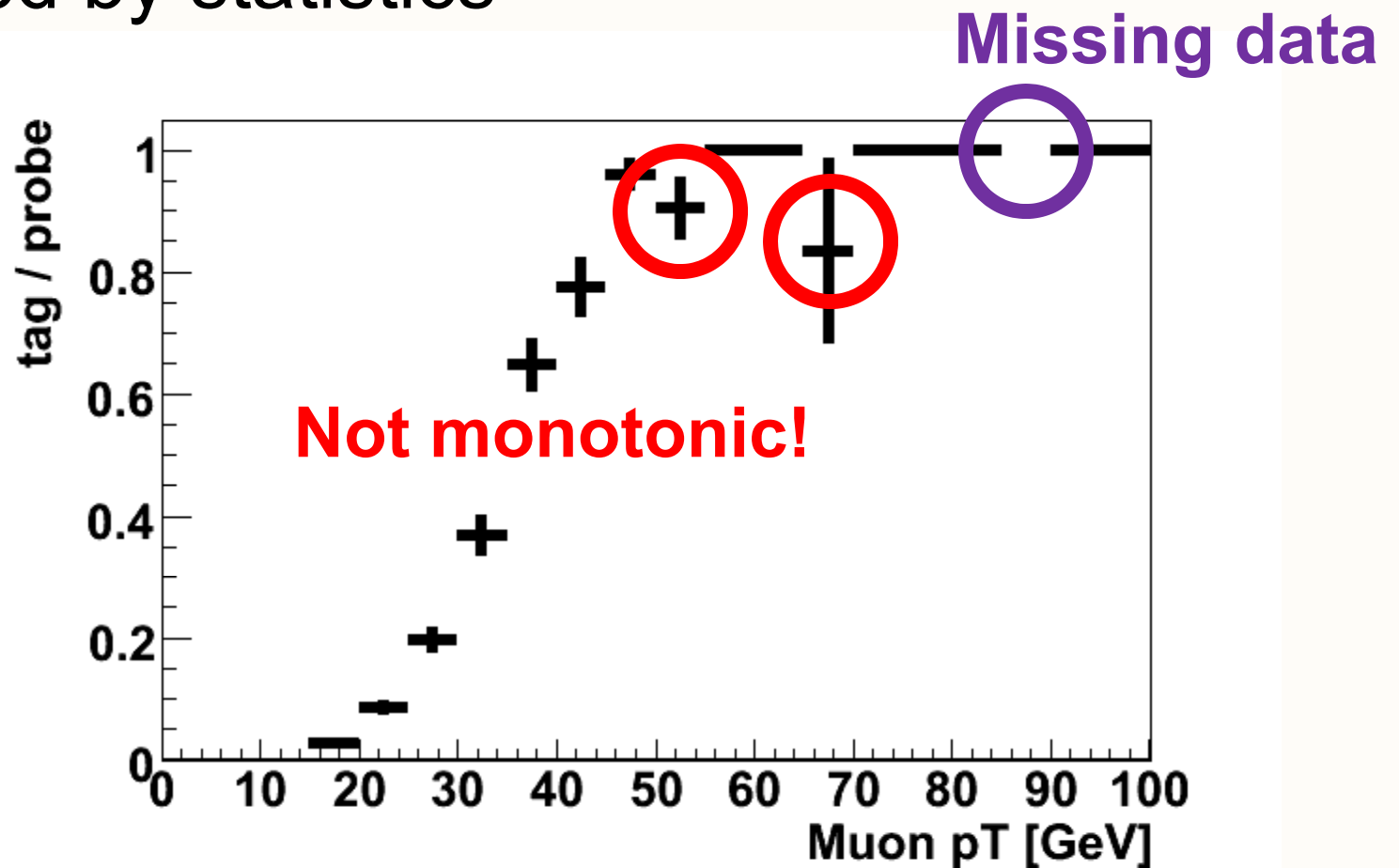


# Statistics, Or: We Know Better!



Sampling "known" distribution

Influenced by statistics

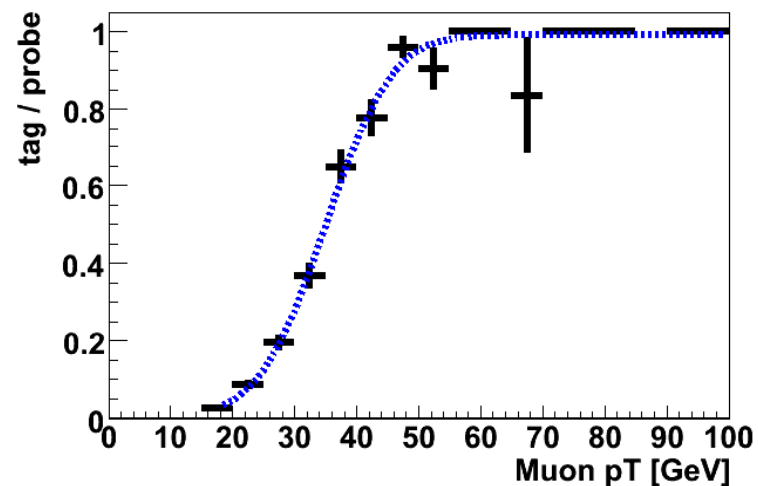
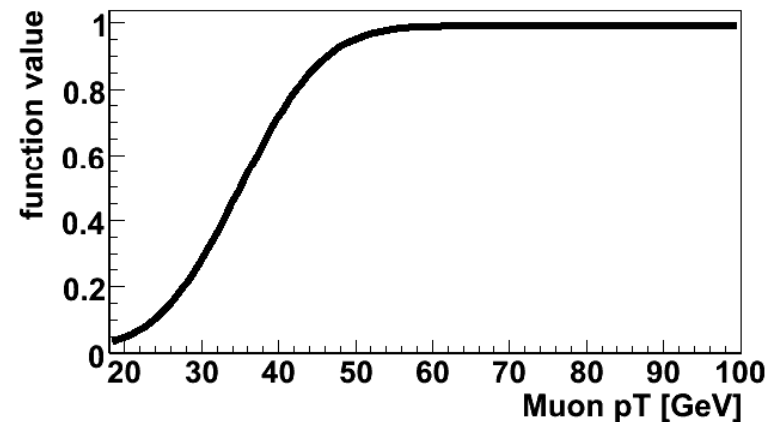


# Fit



Combine our knowledge with statistics / data by fitting a distribution:

1. Find appropriate function with parameters
2. Fit function to distribution



# Fitting: The Function



Finding the proper function involves:

- behavioral analysis:  
starts at 0, goes to constant, monotonic,...
- physics interpretation:  
"E proportional to  $\sin^2(\phi)$ "
- having a good knowledge of typical functions  
(see TMath)
- finding a good compromise between  
generalization ("constant") and precision  
("polynomial 900<sup>th</sup> degree")





# Fitting: The Function

Finding the proper function involves:

- behavioral analysis:  
starts at 0, goes to constant, monotonic, ...
- physics interpretation:  
"E proportional to  $\sin^2(\phi)$ "
- having a good knowledge of typical functions  
(see TMath)
- finding a good compromise between  
generalization ("constant") and precision  
("polynomial 900<sup>th</sup> degree")

**MAGIC**

# Fitting: Parameters

Let's take "erf"  $\text{erf}(x) / 2 + 0.5$

Free parameters:

$$(\text{erf}((x - [0]) / [1]) / 2 + 0.5) * [2]$$

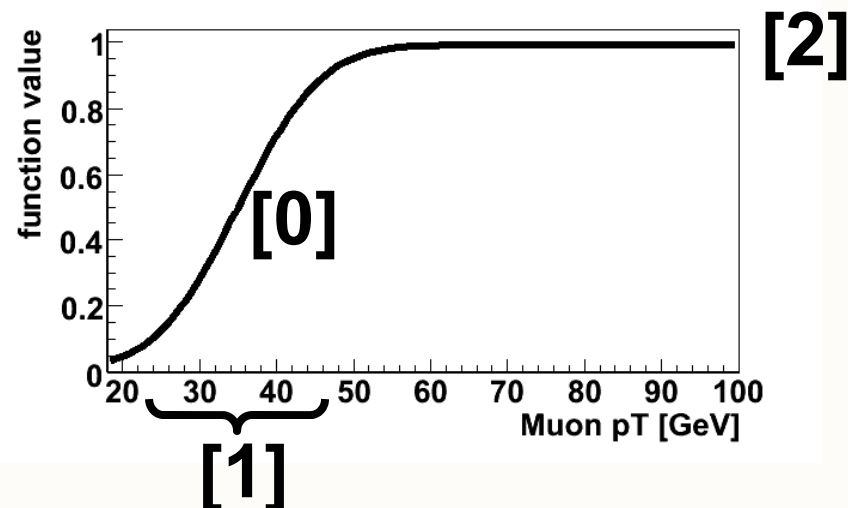
[0]: x @ center of the slope

[1]: width of the slope

[2]: maximum efficiency

How do I know?

Study function behavior

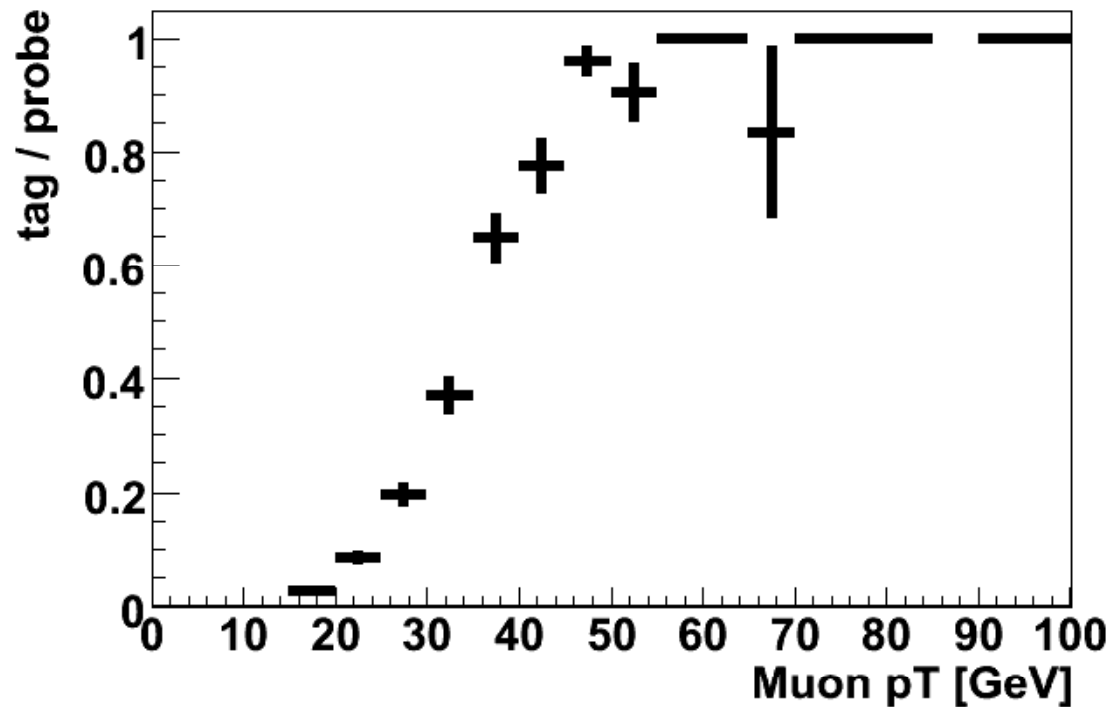




# Fitting: Parameter Init

Need to initialize parameters!

sensible: [0]: 35, [1]: 10, [2]: 1



# Fitting: The Math



Fitting = finding parameters such that

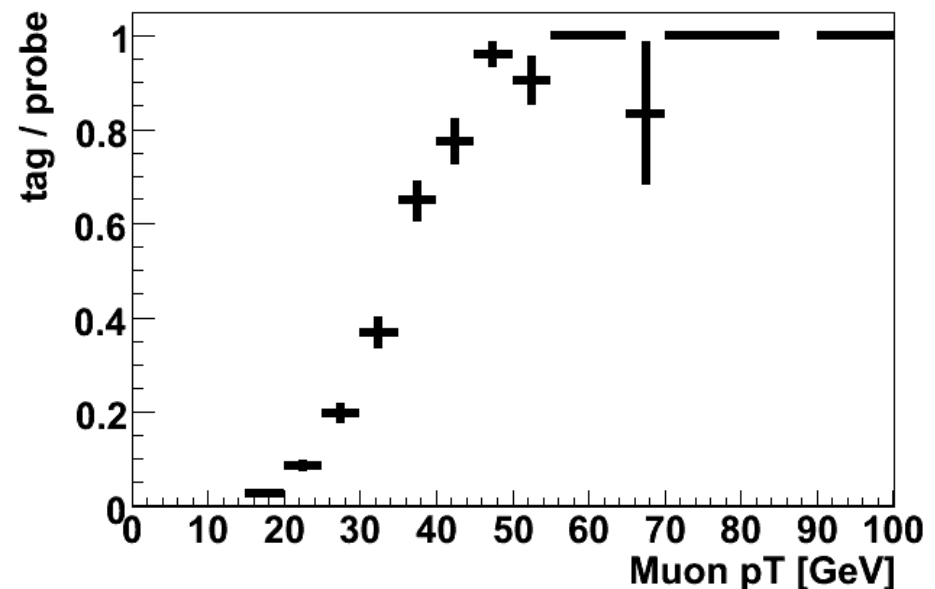
$$f(x) - \text{hist}(x)$$

minimal for all  $x$  [or any similar measure]

Histogram with errors:

$$f(x) - \text{hist}(x) / \text{err}(x)$$

[or similar]



# Fitting In ROOT



Define fit function:

```
TF1* f = new TF1("myfit",  
  "(TMath::Erf((x-[0])/[1])/2.+0.5)*[2]"  
  0., 100.);
```

Set parameters:

```
f->SetParameter(0, 35.);  
f->SetParameter(1, 10.);  
f->SetParameter(2, 1.);
```

Fit it to the histogram:

```
hist->Fit(f);
```

# Fitting Result

Result of fit printed or

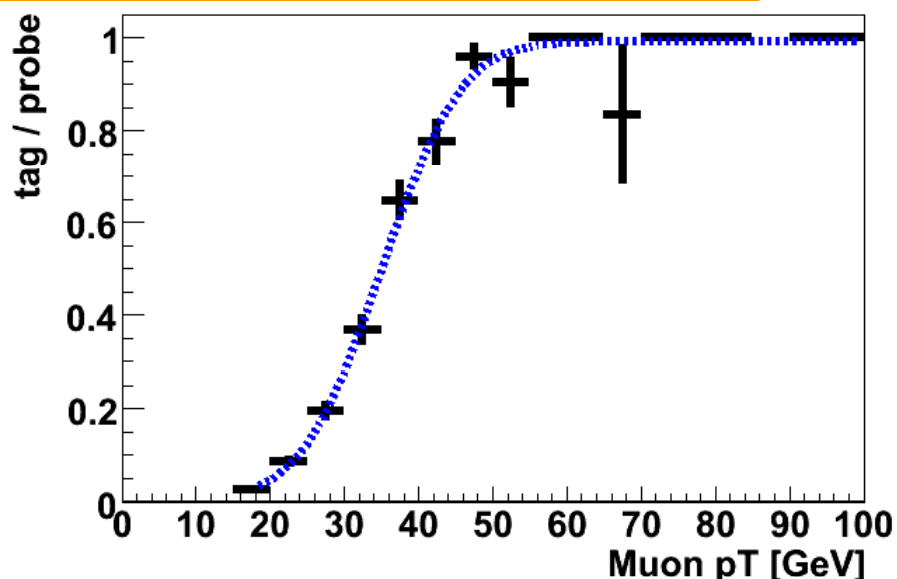
```
cout << f->GetParameter(0) << endl;
```

[0]: 34.9

[1]: 12.1

[2]: 0.98

which means:



```
(TMath::Erf((x-34.9)/12.1)/2.+0.5)*0.98
```

Get efficiency at pT=42GeV:

```
f->Eval(42.)
```

# Analysis: Recap



We started with the trigger problem – and ended with an answer

You now know

- how to determine trigger efficiency from triggered data
- why large samples are relevant
- what fitting is, how it works, when to do it, and how it's done with ROOT.



# Interactive Data Analysis with PROOF

Bleeding Edge Physics with  
Bleeding Edge Computing



# Parallel Analysis: PROOF



Some numbers (from Alice experiment)

- 1.5 PB ( $1.5 * 10^{15}$ ) of raw data per year
- 360 TB of ESD+AOD\* per year (20% of raw)
- One pass at 15 MB/s will take 9 months!

Parallelism is the only way out!

\* ESD: Event Summary Data      AOD: Analysis Object Data



# PROOF

Huge amounts of events, hundreds of CPUs

Split the job into  $N$  events / CPU!

PROOF for TSelector based analysis:

- **start** analysis locally ("client"),
- PROOF **distributes** data and code,
- lets CPUs ("workers") **run** the analysis,
- **collects** and combines (merges) data,
- shows analysis **results** locally
- More **dynamic** than a batch system

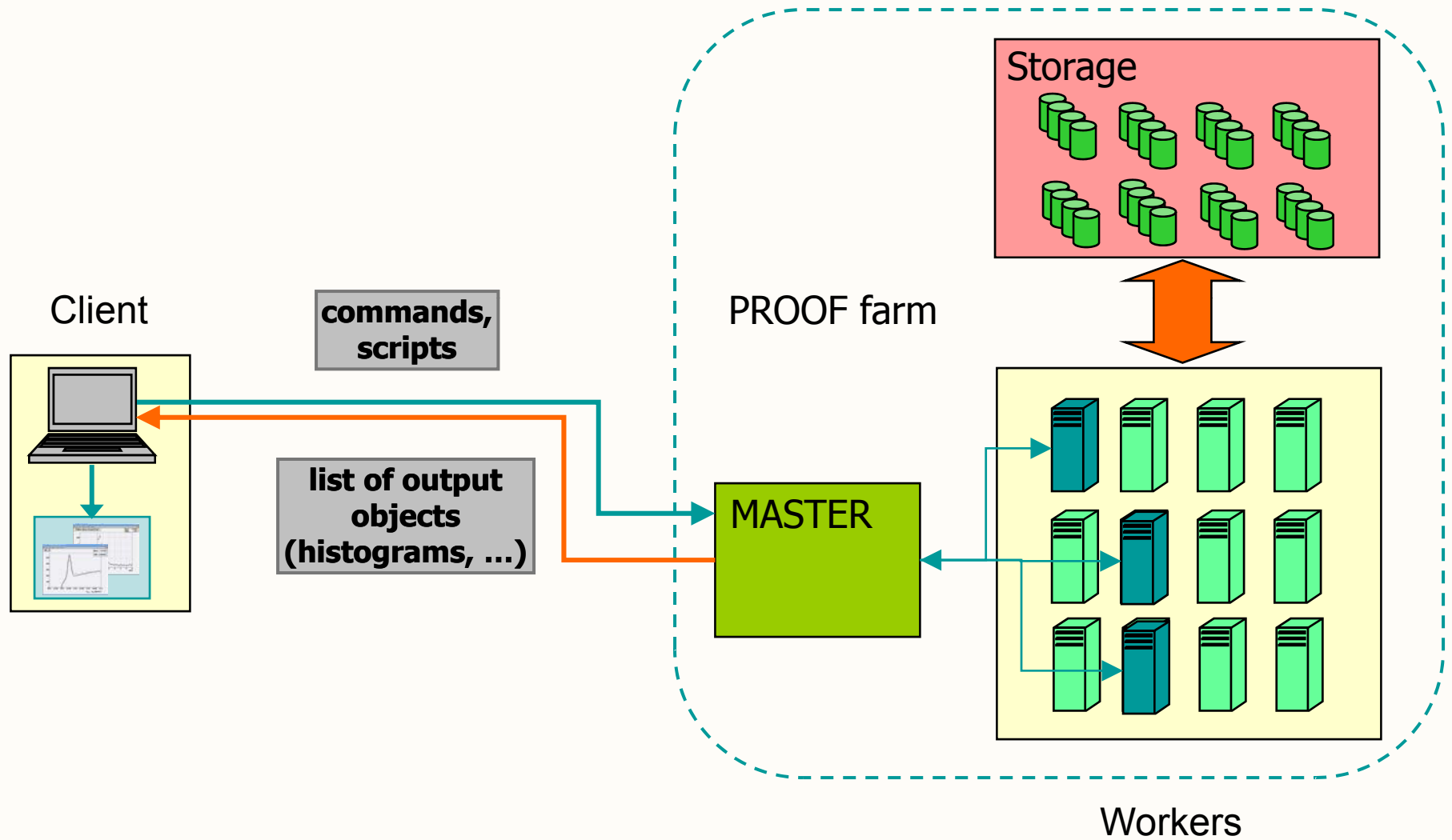
Including on-the-fly status reports!

# Interactive!



- Start analysis
- Watch status while running
- Forgot to create a histogram?
  - Interrupt the process
  - Modify the selector
  - Re-start the analysis

# PROOF



# Creating a session



To create a PROOF session from the ROOT prompt, just type:

```
TProof *p = TProof::Open("master")
```

where "master" is the hostname of the master machine on the PROOF cluster

# PROOF Analysis



- Example of local TChain analysis

```
// Create a chain of trees
root[0] TChain *c = new TChain("myTree");
root[1] c->Add("http://www.any.where/file1.root");
root[2] c->Add("http://www.any.where/file2.root");

// MySelector is a TSelector
root[3] c->Process("MySelector.C+");
```

# PROOF Analysis



- Same example with PROOF

```
// Create a chain of trees
root[0] TChain *c = new TChain("myTree");
root[1] c->Add("http://www.any.where/file1.root");
root[2] c->Add("http://www.any.where/file2.root");

// Start PROOF and tell the chain to use it
root[3] TProof::Open("masterURL");
root[4] c->SetProof();

// Process goes via PROOF
root[5] c->Process("MySelector.C+");
```

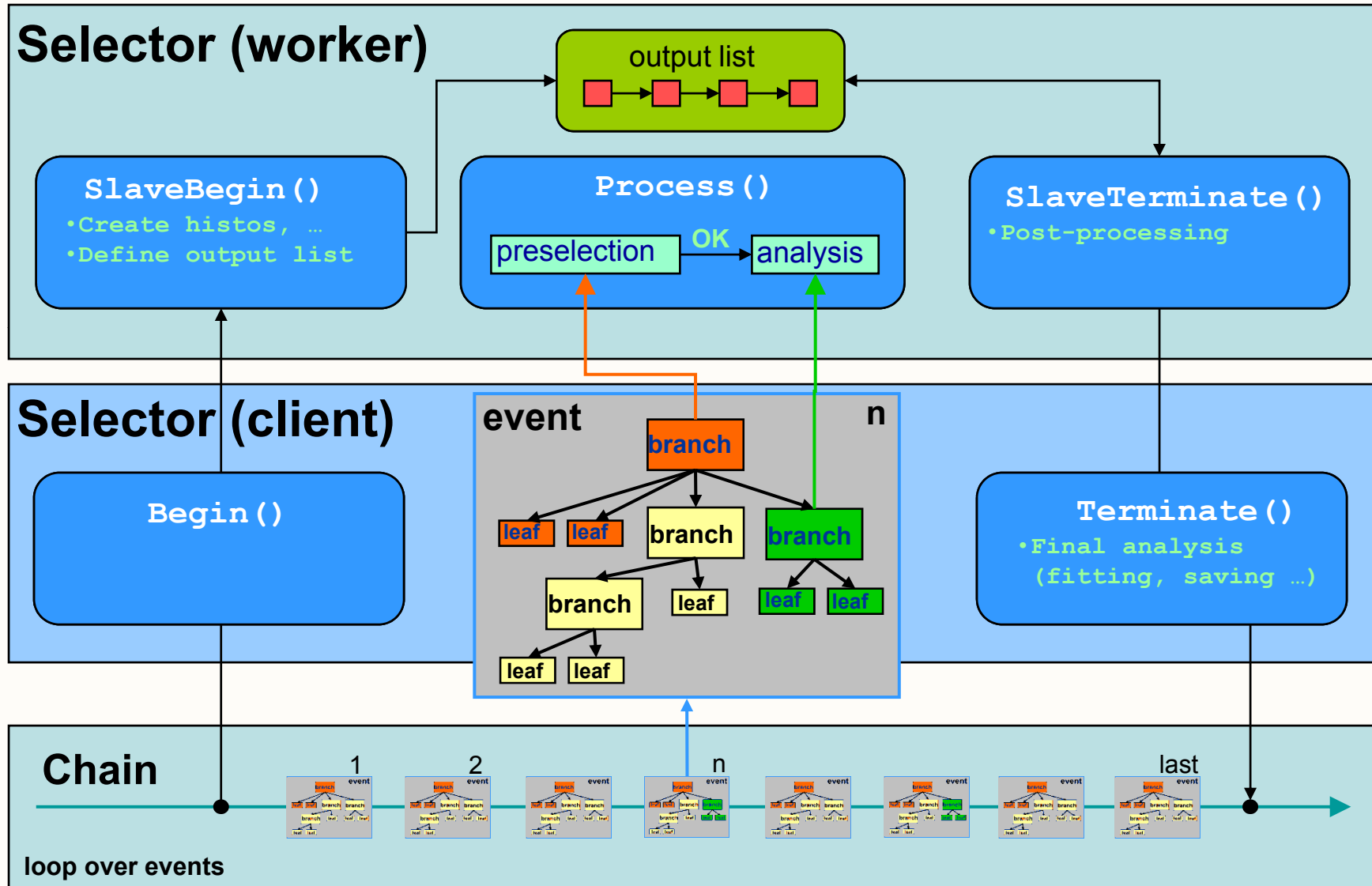
# TSelector & PROOF



- **Begin()** called on the **client** only
- **SlaveBegin()** called on each **worker**: create histograms
- **SlaveTerminate()** rarely used; post processing of partial results before they are sent to master and merged
- **Terminate()** runs on the **client**: save results, display histograms, ...



# PROOF Analysis



# Output List (result of the query)



- Each worker has a partial output list
- Objects have to be added to the list in `TSelector::SlaveBegin()` e.g.:

```
fHist = new TH1F("h1", "h1", 100, -3., 3.);  
fOutput->Add(fHist);
```

- At the end of processing the output list gets sent to the master
- The Master merges objects and returns them to the client. Merging is e.g. "Add()" for histograms, appending for lists and trees

# Results



At the end of `Process()`, the output list is accessible via `gProof->GetOutputList()`

```
// Get the output list
root[0] TList *output = gProof->GetOutputList();
// Retrieve 2D histogram "h2"
root[1] TH2F *h2 = (TH2F*)output->FindObject("h2");
// Display the histogram
root[2] h2->Draw();
```

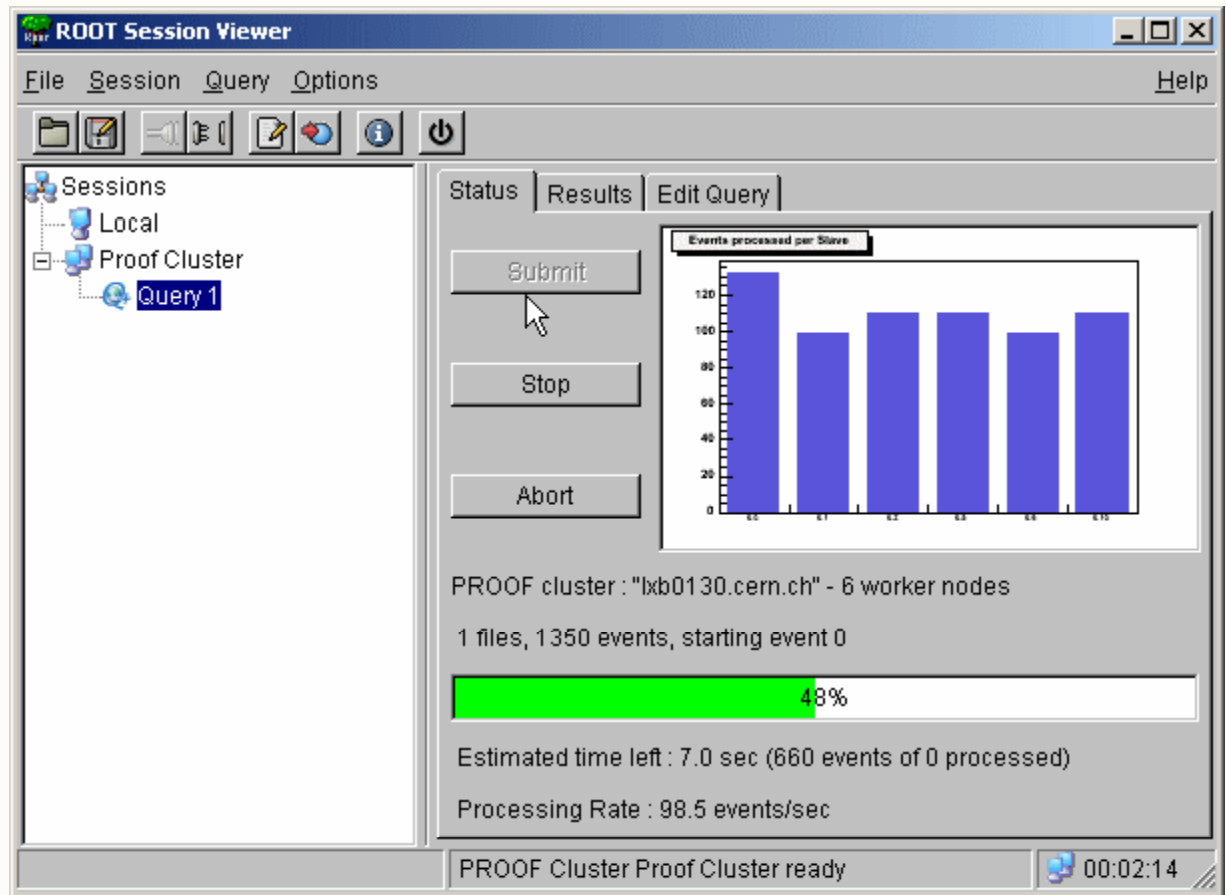
# PROOF GUI Session



Starting a PROOF GUI session is trivial:

```
TProof::Open()
```

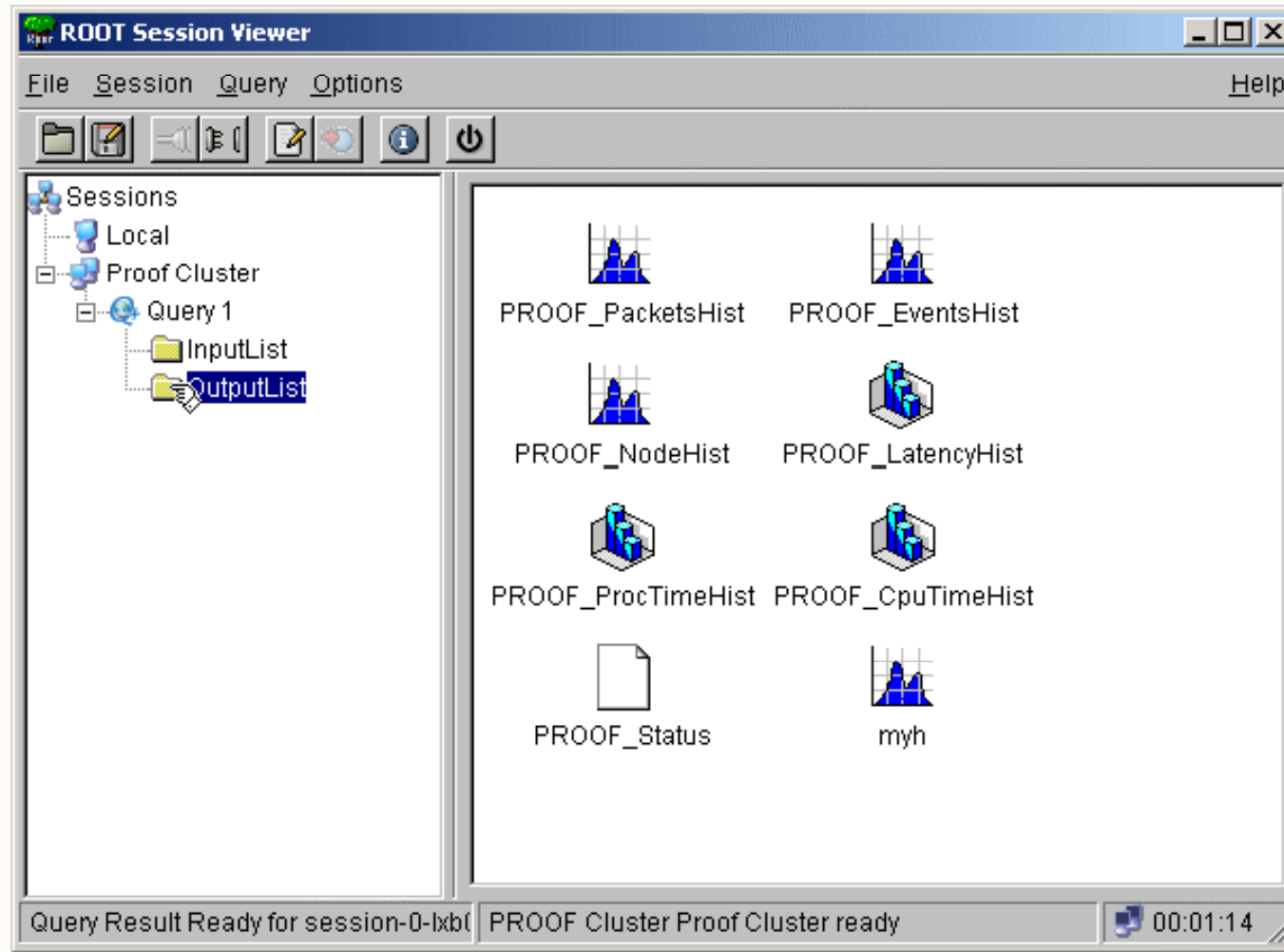
Opens GUI:



# PROOF GUI Session – Results



Results accessible via TSessionViewer, too:



# PROOF Documentation

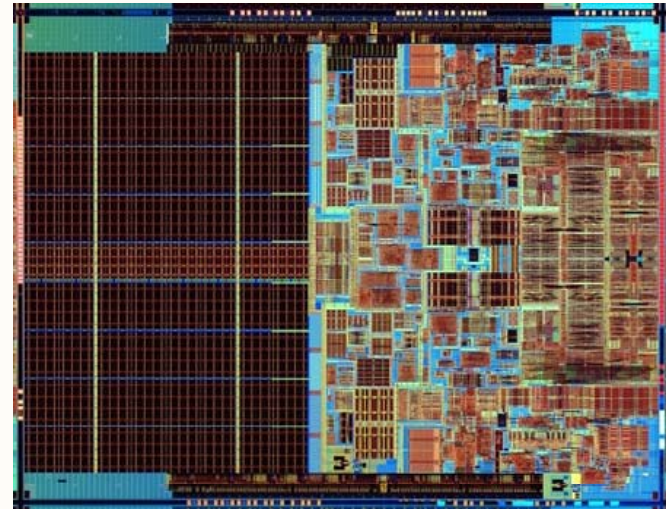


Documentation available online at

[root.cern.ch/twiki/bin/view/ROOT/PROOF](http://root.cern.ch/twiki/bin/view/ROOT/PROOF)

But of course you need a little cluster of CPUs

Like your multicore laptop!



# Summary



You've learned:

- analyzing a TTree can be easy and efficient
- integral part of physics is counting
- ROOT provides histogramming and fitting
- > 1 CPU: use PROOF!

Looking forward to hearing from you:

- as a user (help! bug! suggestion!)
- and as a developer!